

Java Concurrency Live(li)ness Lab

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Live(li)ness

- **Liveness (labs 1-4)**
 - No deadlocks, no livelocks
- **Liveliness (labs 5.1-5.5)**
 - Full CPU utilization, no contention



tinyurl.com/jfokus2017

Avoiding Liveness Hazards



Avoiding Liveness Hazards

- **Fixing safety problems can cause liveness problems**
 - **Don't indiscriminately sprinkle "synchronized" into your code**

Deadly Embrace

- **Lock-ordering deadlocks**
 - Typically when you lock two locks in different orders
 - Requires global analysis to make sure your order is consistent
 - Lesson: only ever hold a single lock per thread!

Thread Deadlocks in BLOCKED

- **A deadly embrace amongst synchronized leaves no way of recovery**
 - **We have to restart the JVM**

Lab 1: Deadlock Resolution by Global Ordering



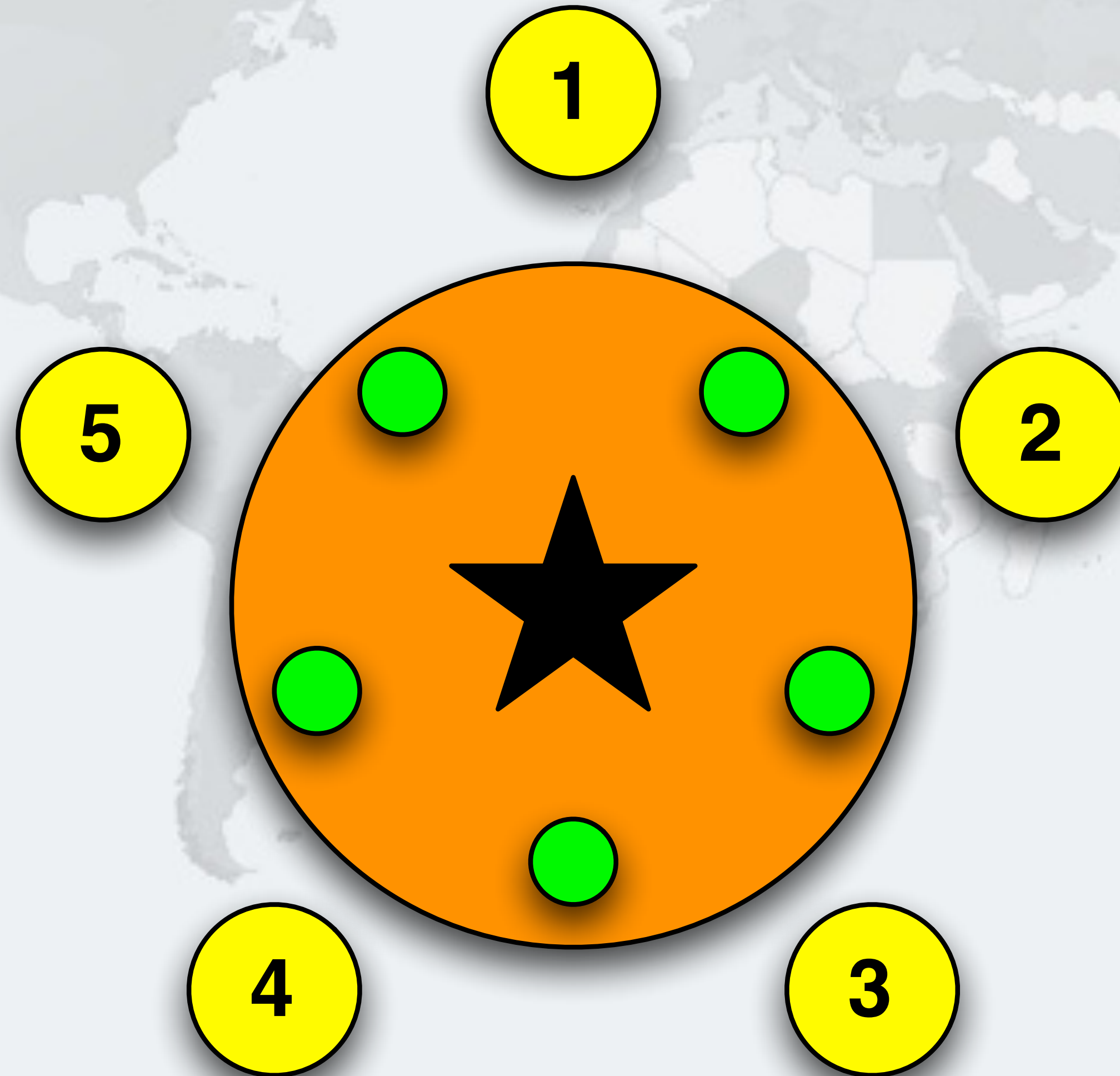
Lab 1: Deadlock resolution by global ordering

- **Classic problem is that of the "dining philosophers"**
 - We changed that to the "drinking philosophers"
 - That is where the word "symposium" comes from
 - sym - together, such as "symphony"
 - poto - drink
 - Ancient Greek philosophers used to get together to drink & think
- **In our example, a philosopher needs two glasses to drink**
 - First he takes the right one, then the left one
 - When he finishes drinking, he returns them and carries on thinking

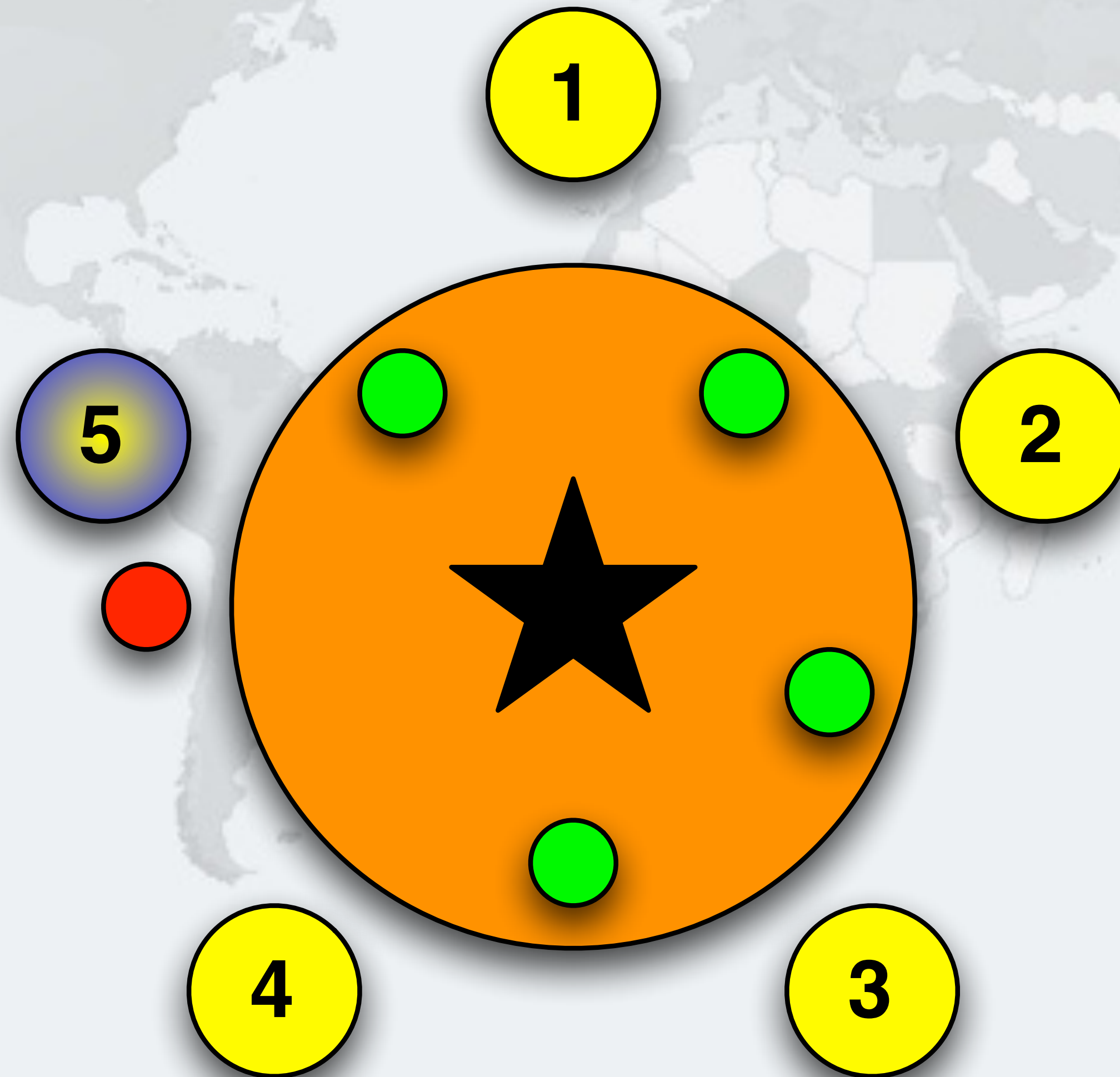
Our Drinking Philosophers

- **Our philosopher needs two glasses to drink**
 - **First he takes the right one, then the left one**
 - **When he's done, he returns the left and then the right**
 - **returns them and carries on thinking**

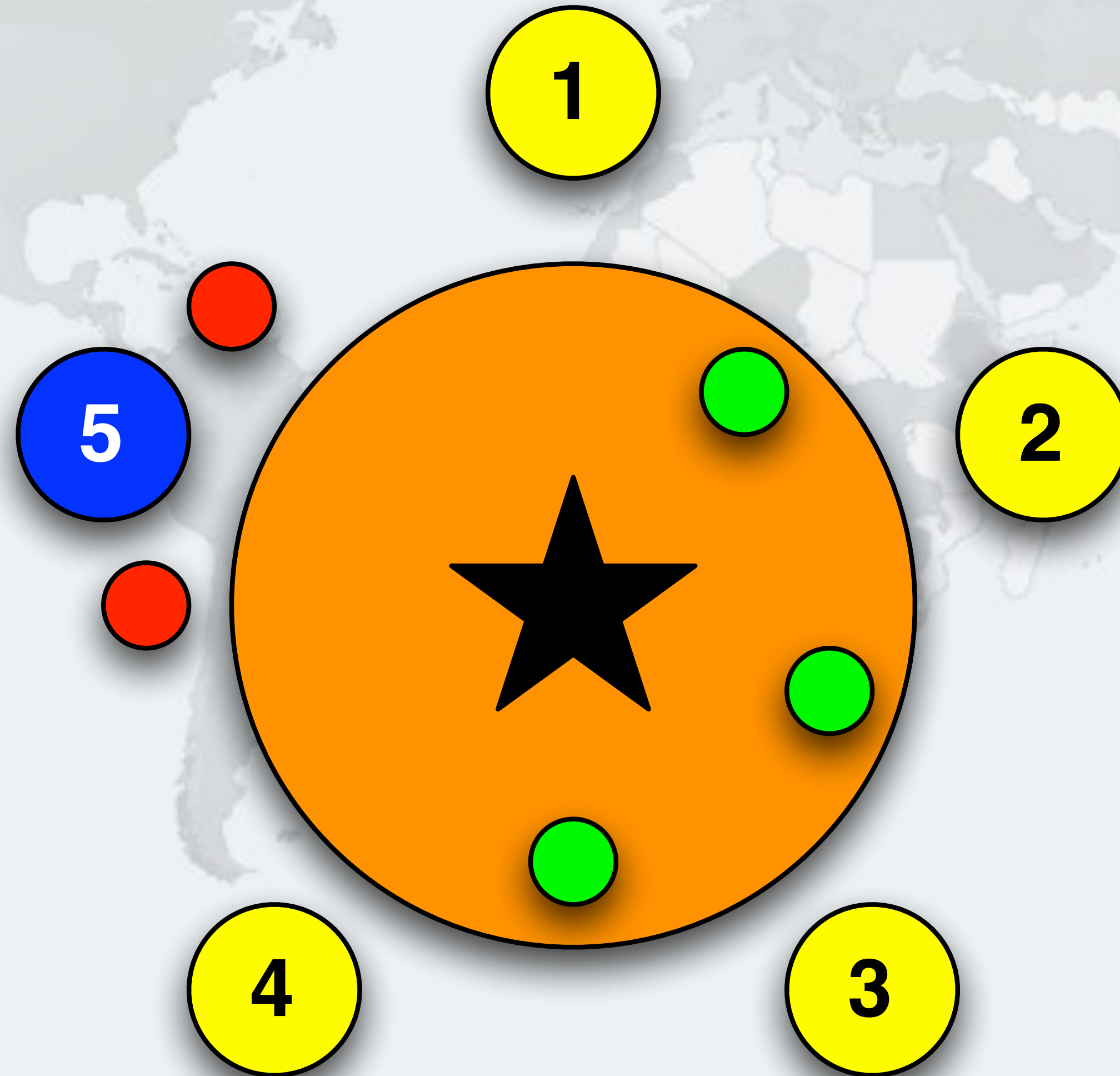
Table is ready, all philosophers are thinking



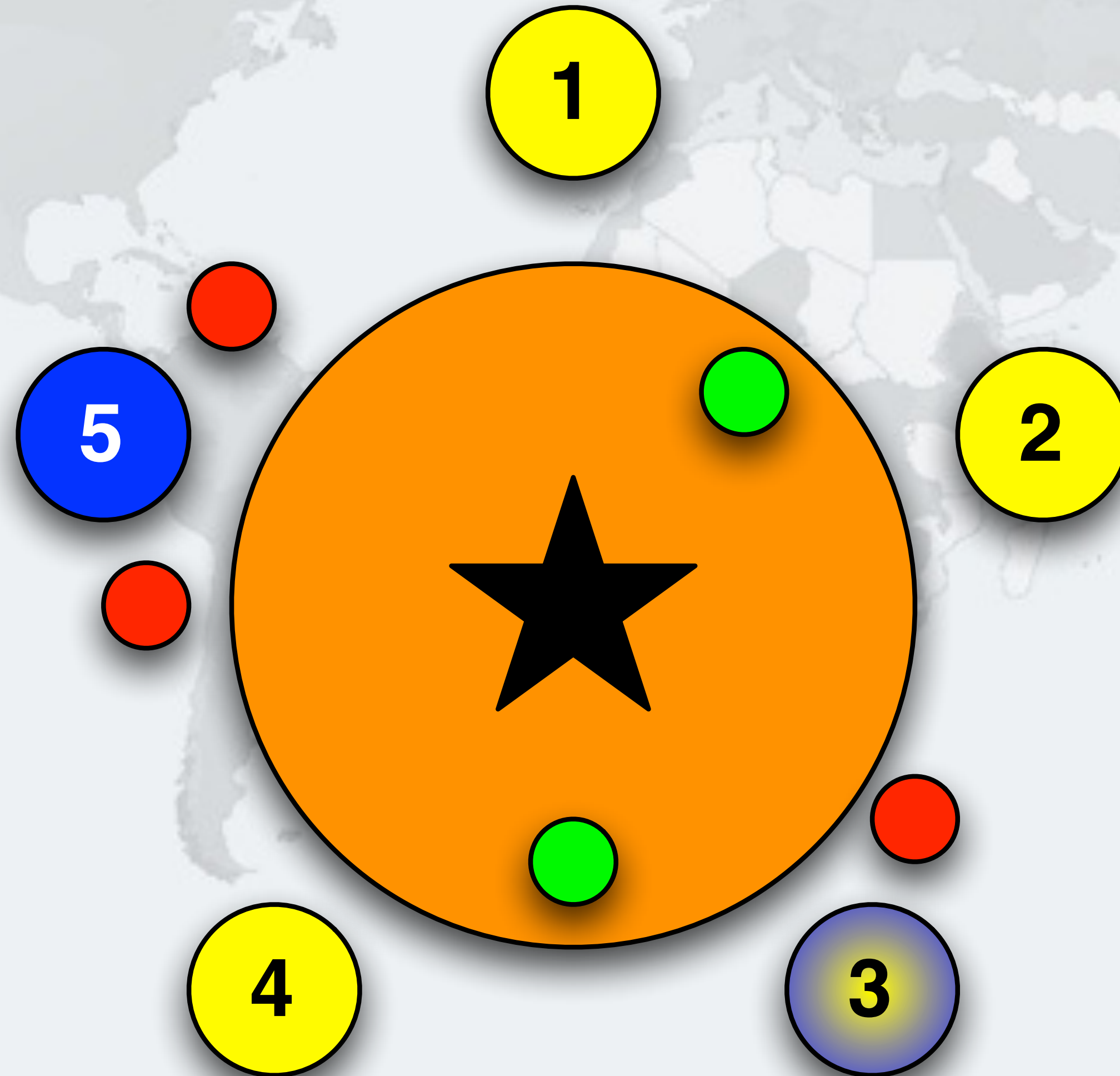
Philosopher 5 wants to drink, takes right cup



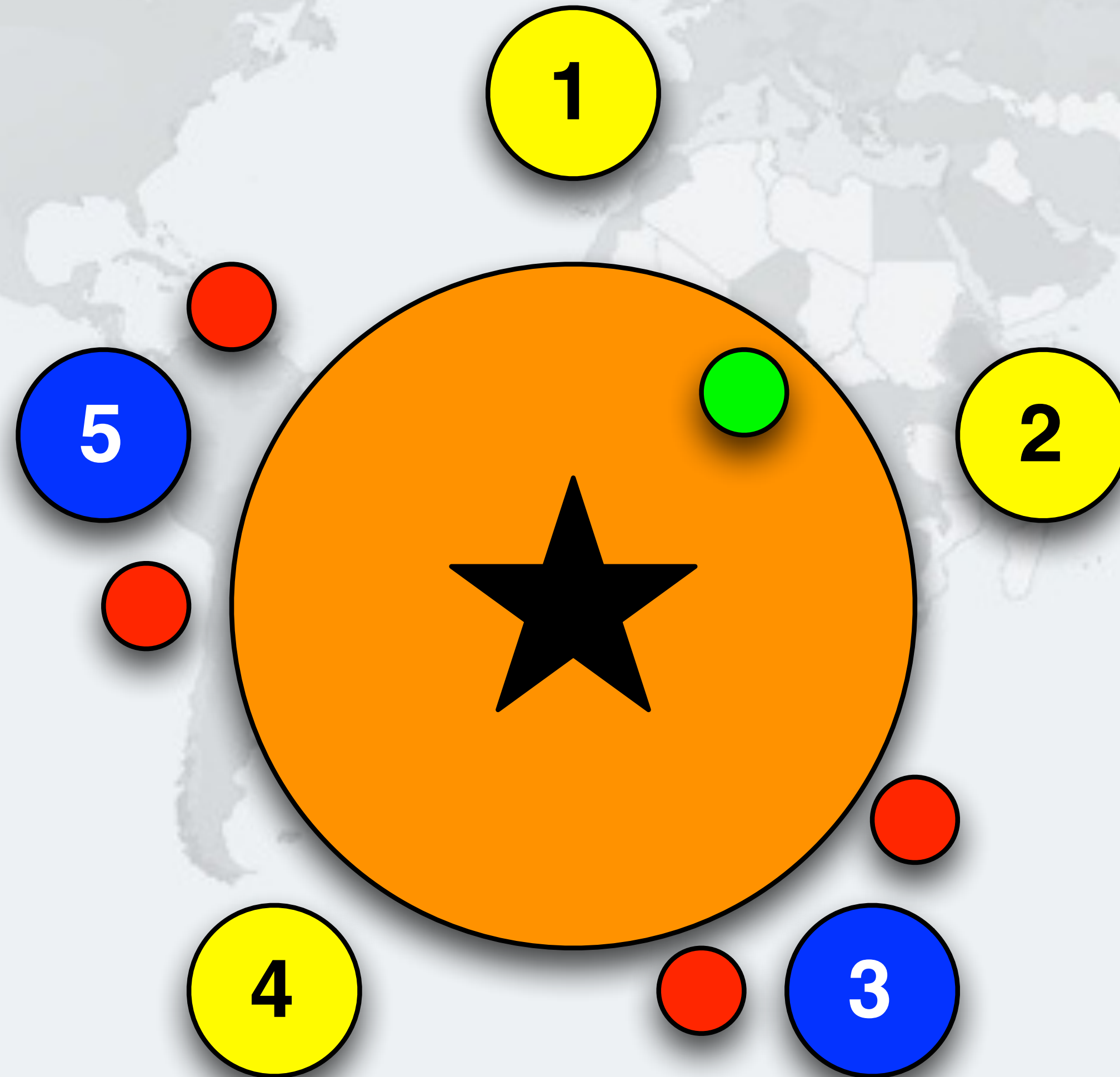
Philosopher 5 is now drinking with both cups



Philosopher 3 wants to drink, takes right cup

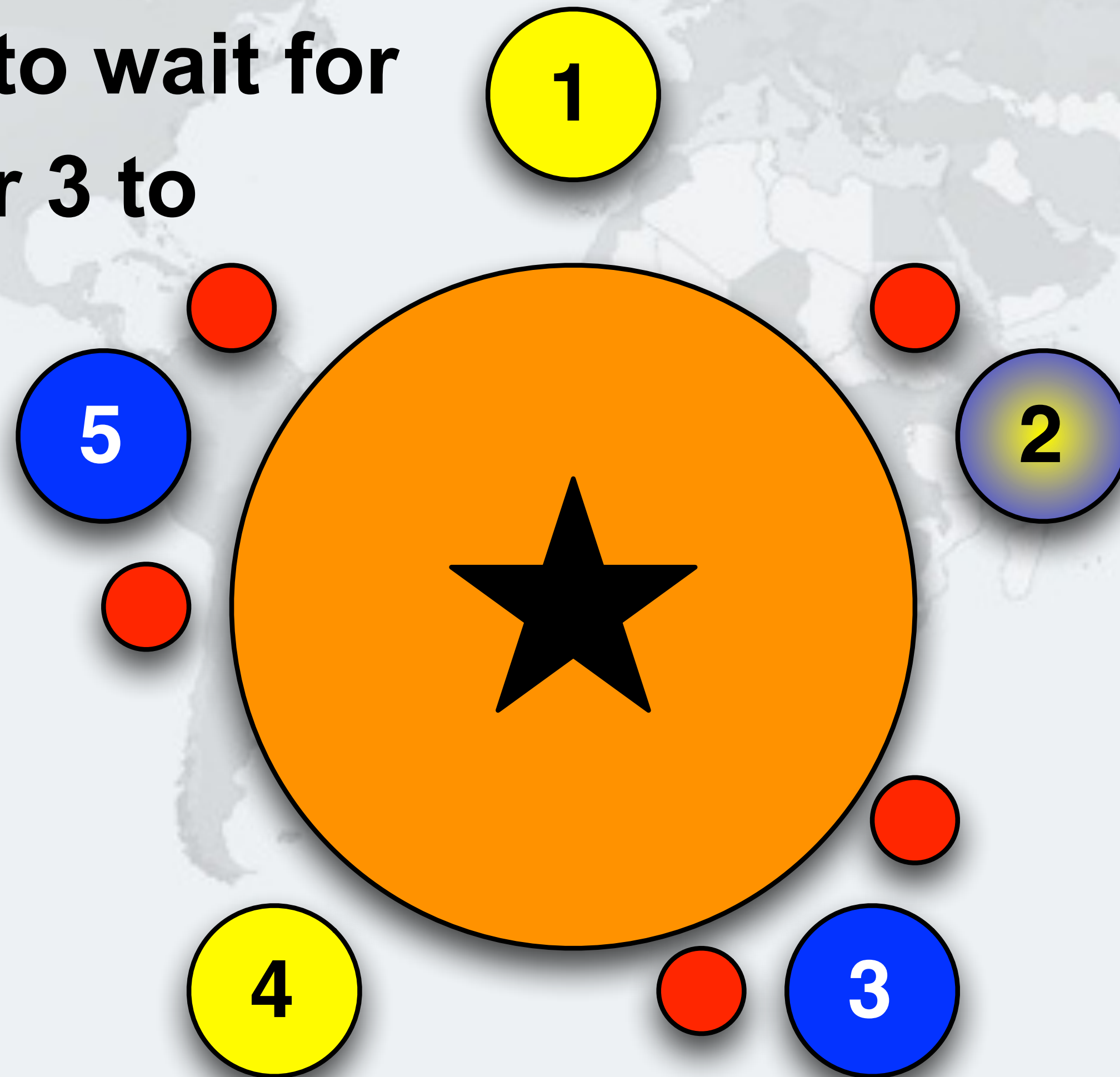


Philosopher 3 is now drinking with both cups

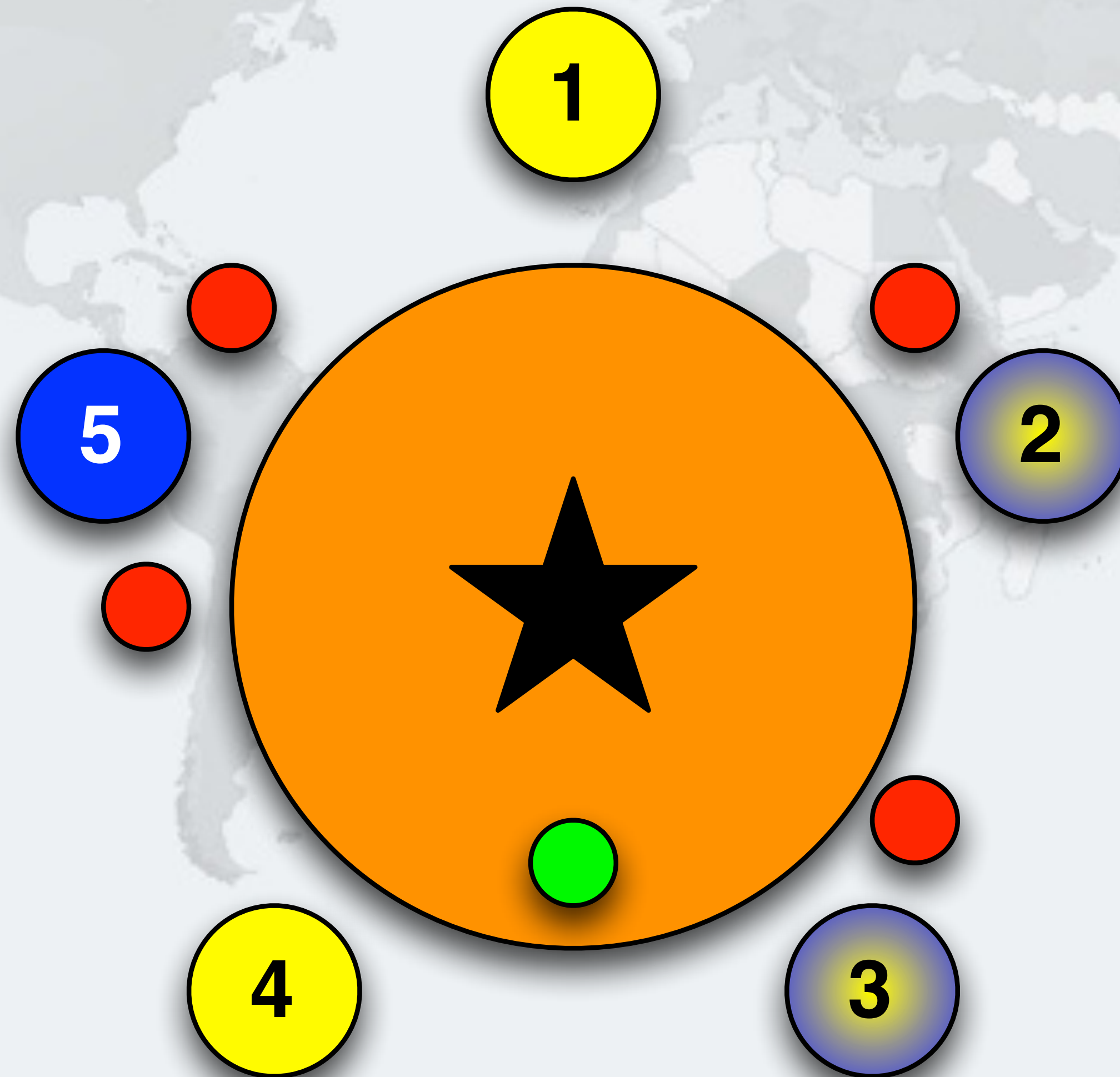


Philosopher 2 wants to drink, takes right cup

- **But he has to wait for Philosopher 3 to finish his drinking session**



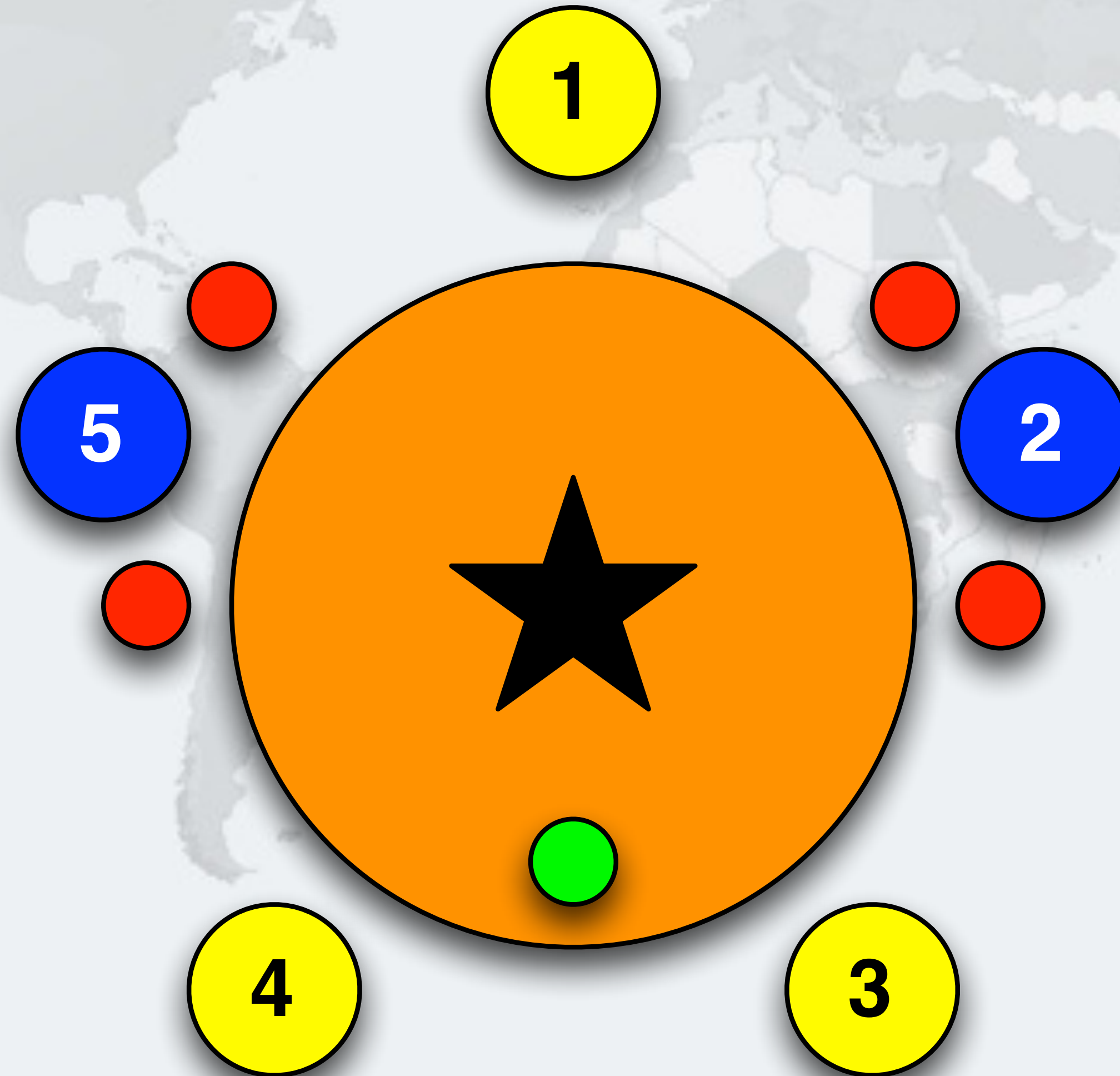
Philosopher 3 finished drinking,
returns left cup



Philosopher 3 returns right cup



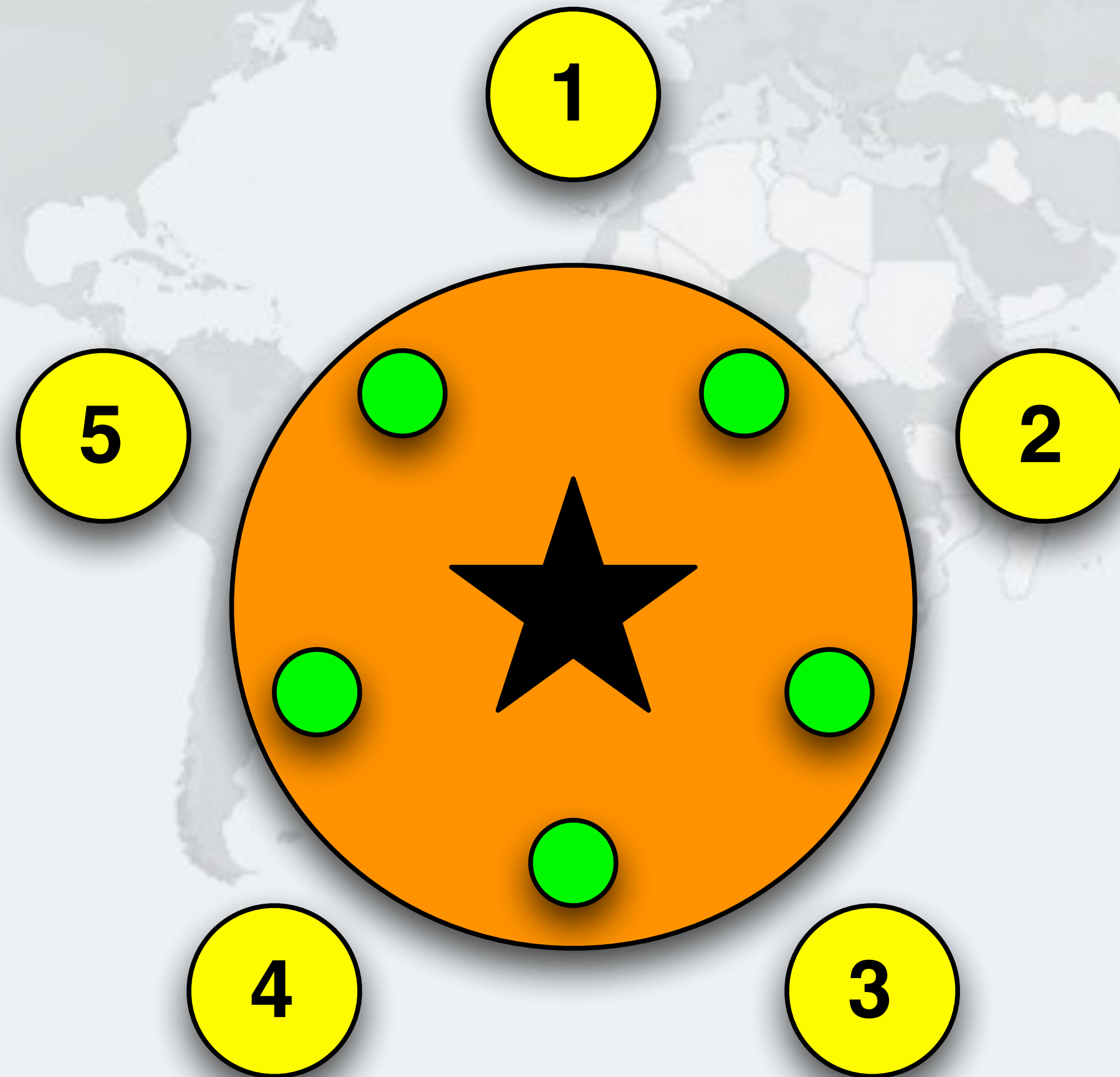
Philosopher 2 is now drinking with both cups



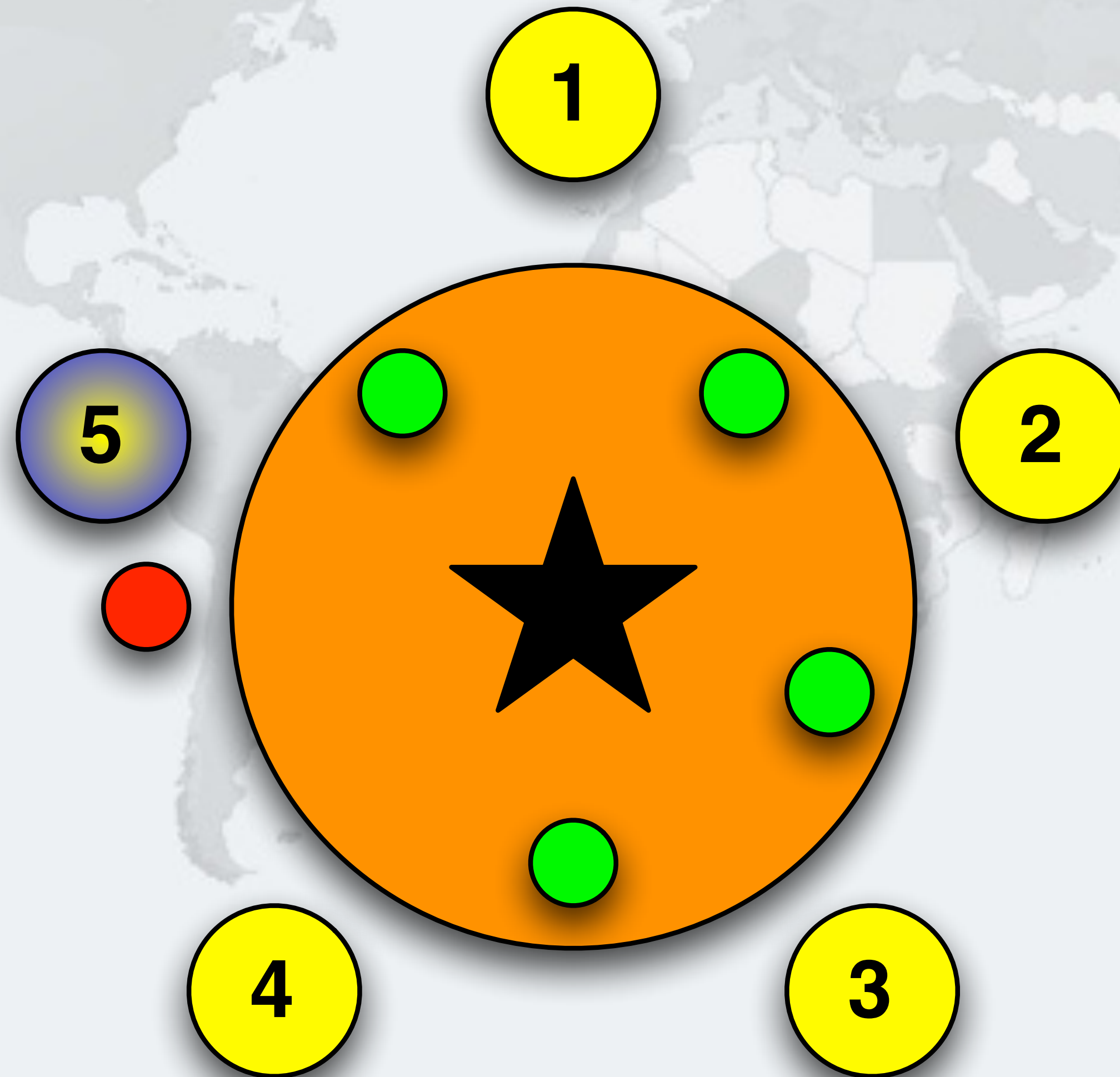
Drinking Philosophers in Limbo

- **The standard rule is that every philosopher first picks up the right cup, then the left**
 - **If all of the philosophers want to drink and they all pick up the right cup, then they all are holding one cup but cannot get the left cup**

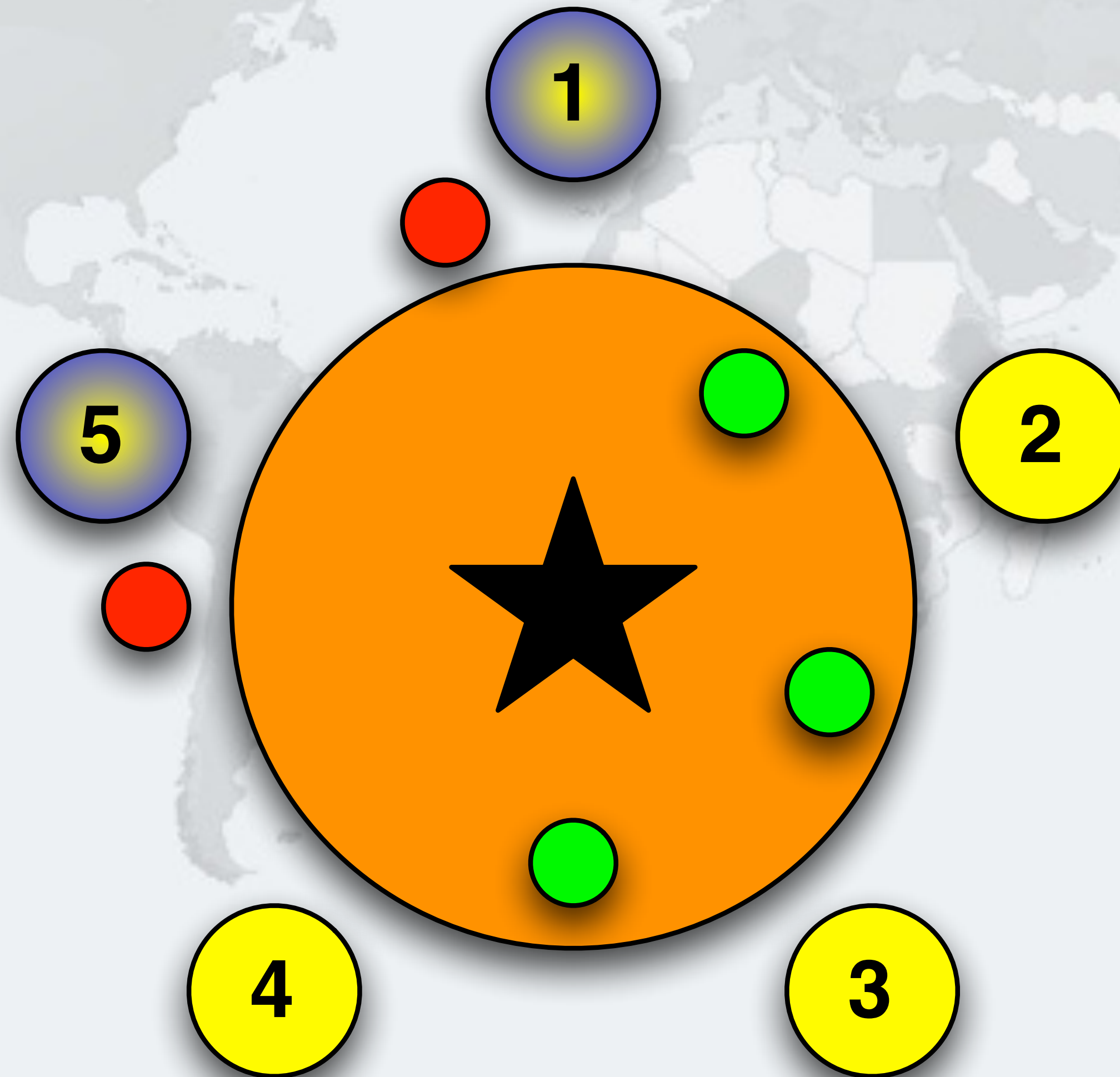
A deadlock can easily happen with this design



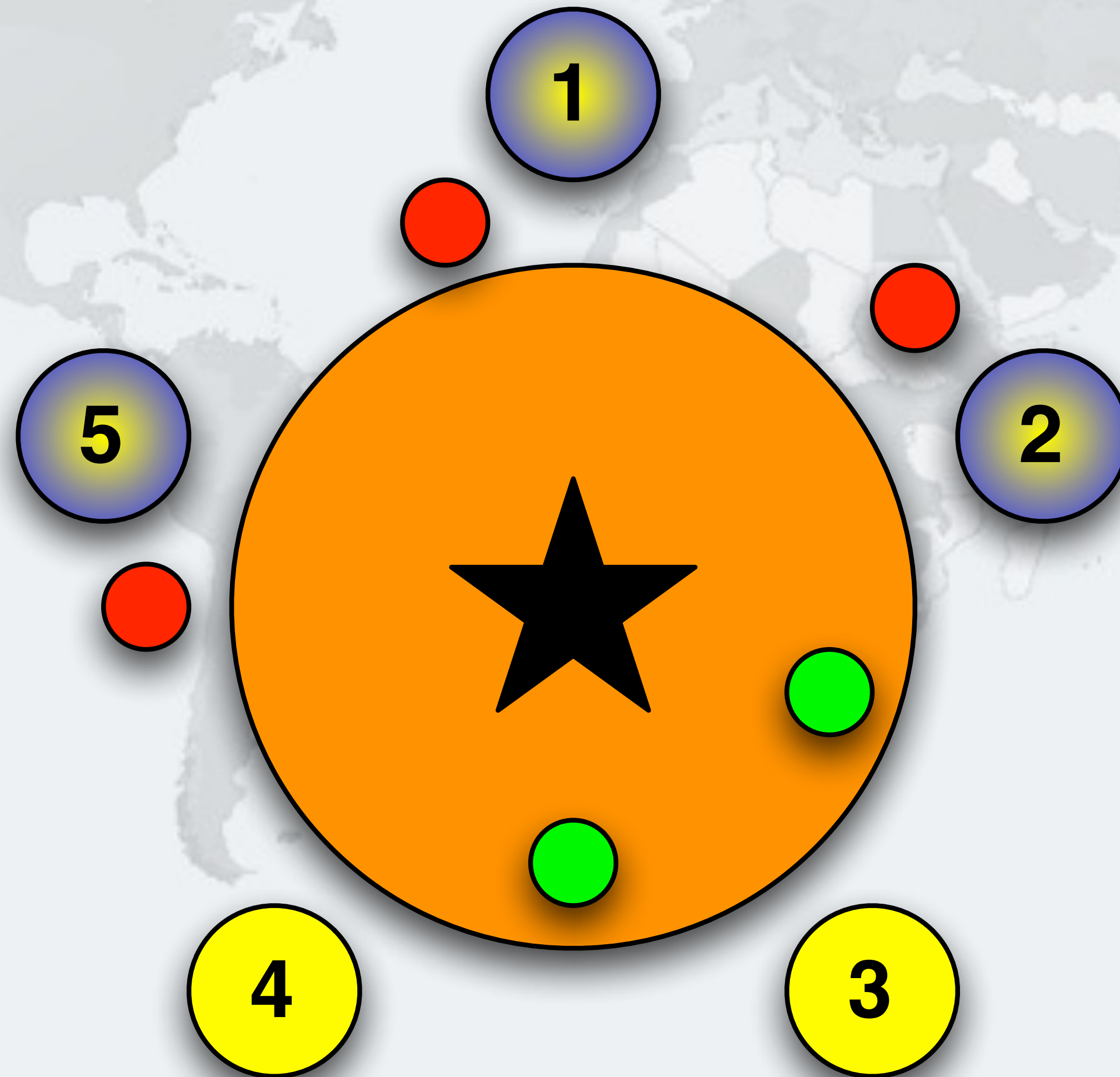
Philosopher 5 wants to drink, takes right cup



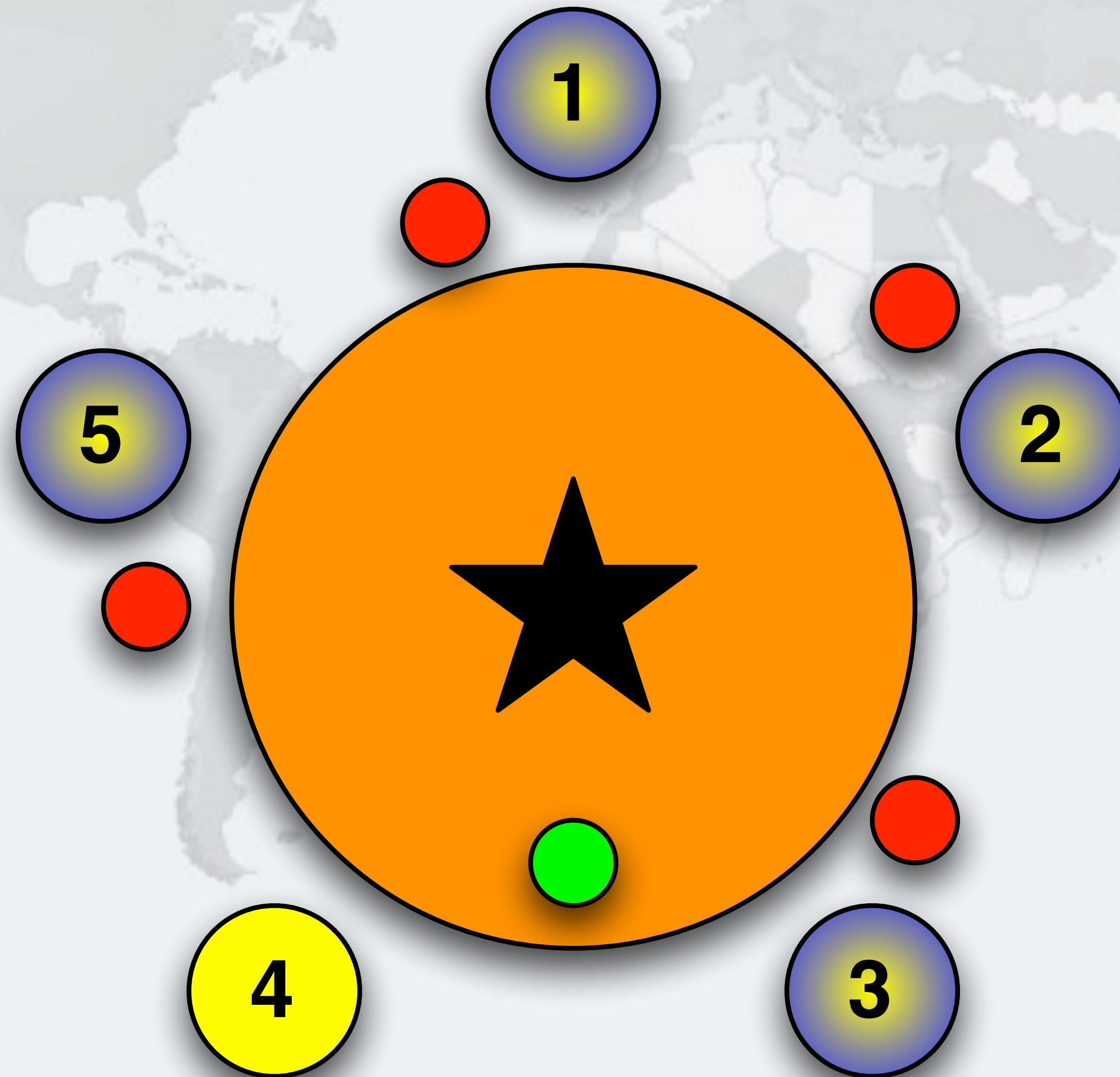
Philosopher 1 wants to drink, takes right cup



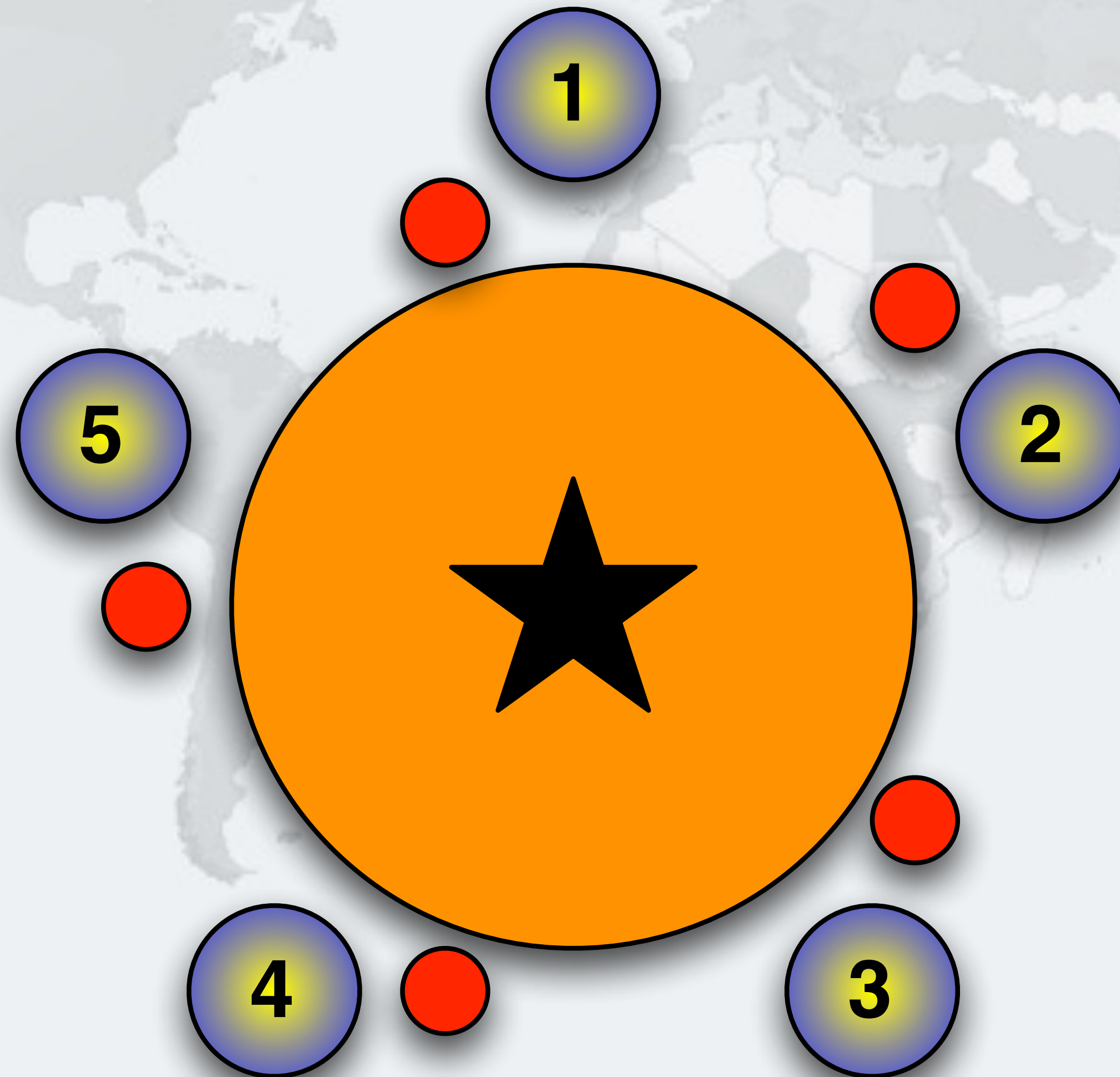
Philosopher 2 wants to drink, takes right cup



Philosopher 3 wants to drink, takes right cup

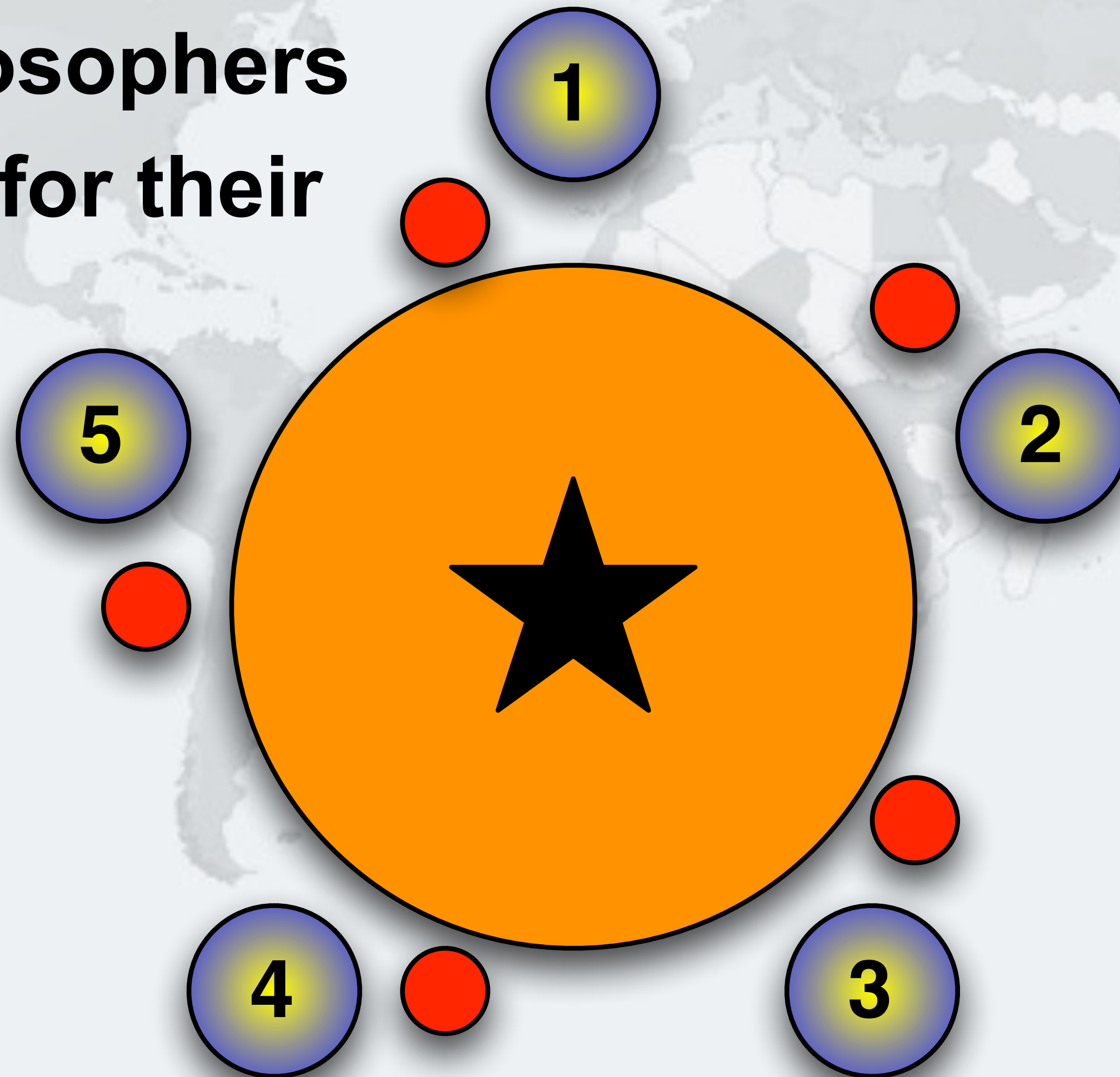


Philosopher 4 wants to drink, takes right cup



Deadlock!

- **All the philosophers are waiting for their left cups, but they will never become available**



Global order with boozing philosophers

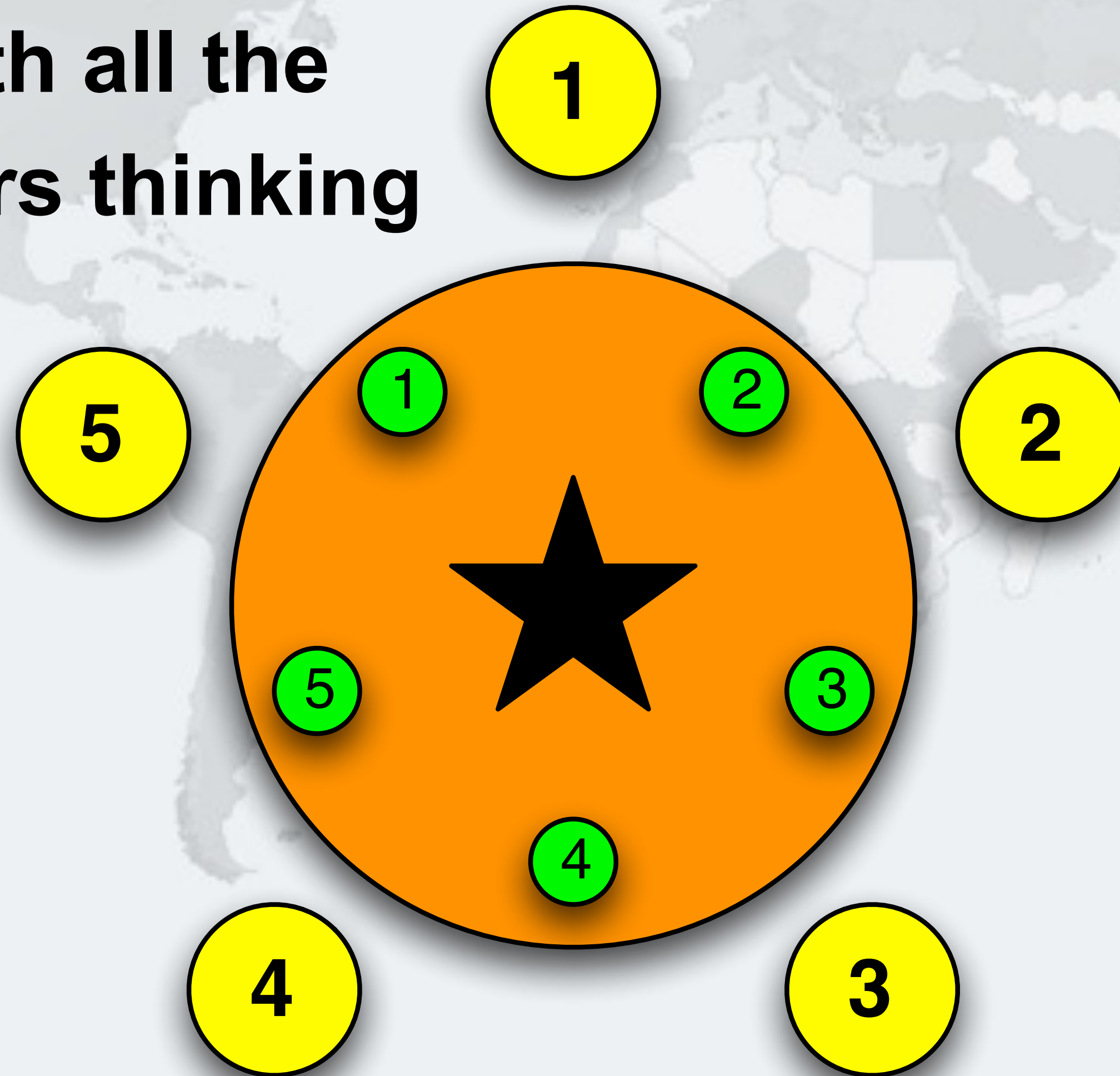
- **If all philosophers hold one cup, we deadlock**
 - **In our solution, we have to prevent that from happening**

fixed order of lock acquisition

- **We can solve the deadlock with the "dining philosophers" by requiring that locks are always acquired in a set order**
 - **For example, we can make a rule that philosophers always first take the cup with the largest number**
 - **If it is not available, we block until it becomes available**
 - **And return the cup with the lowest number first**

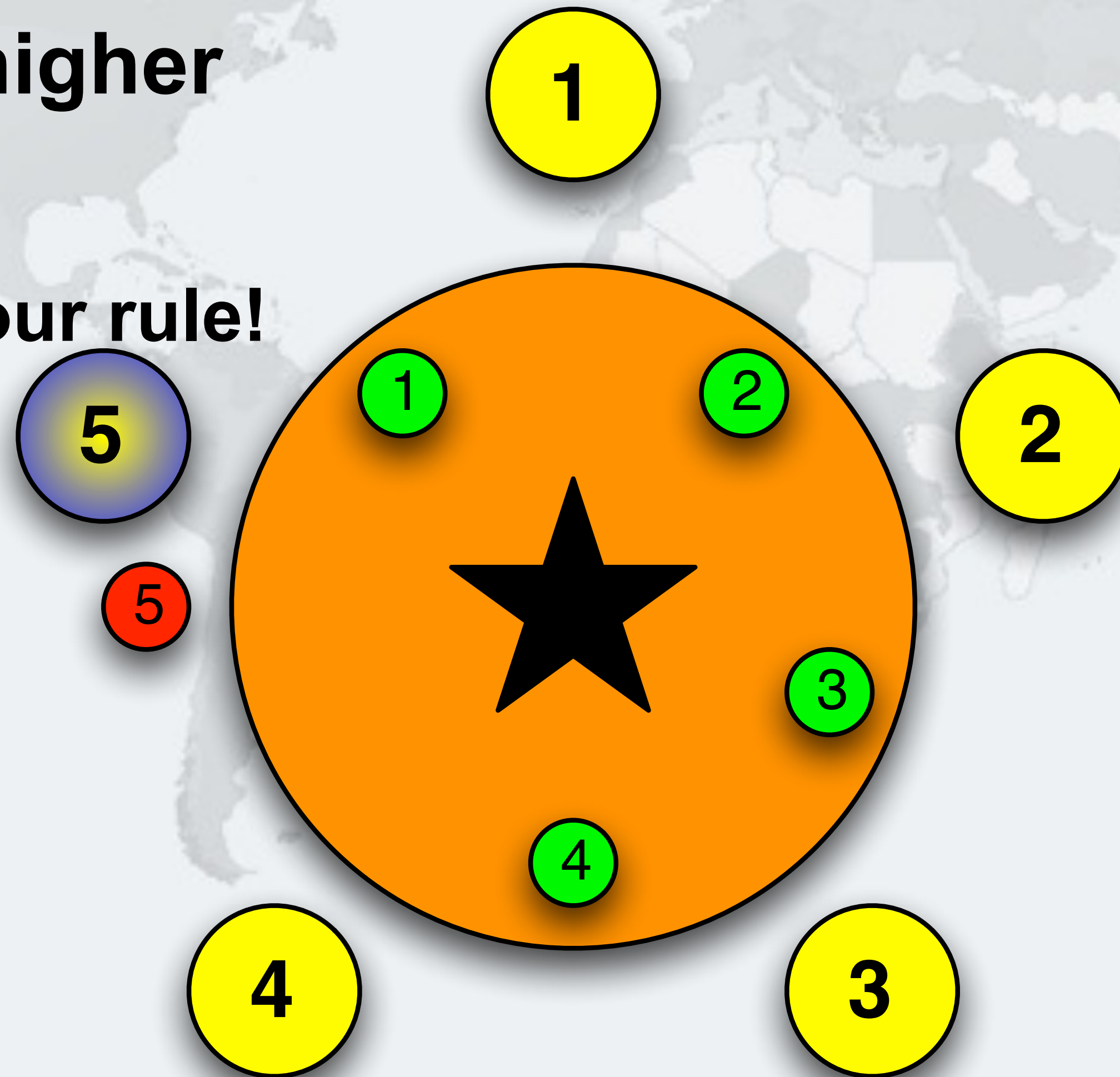
Global Lock ordering

- **We start with all the philosophers thinking**



Philosopher 5 takes cup 5

- **Cup 5 has higher number**
 - Remember our rule!

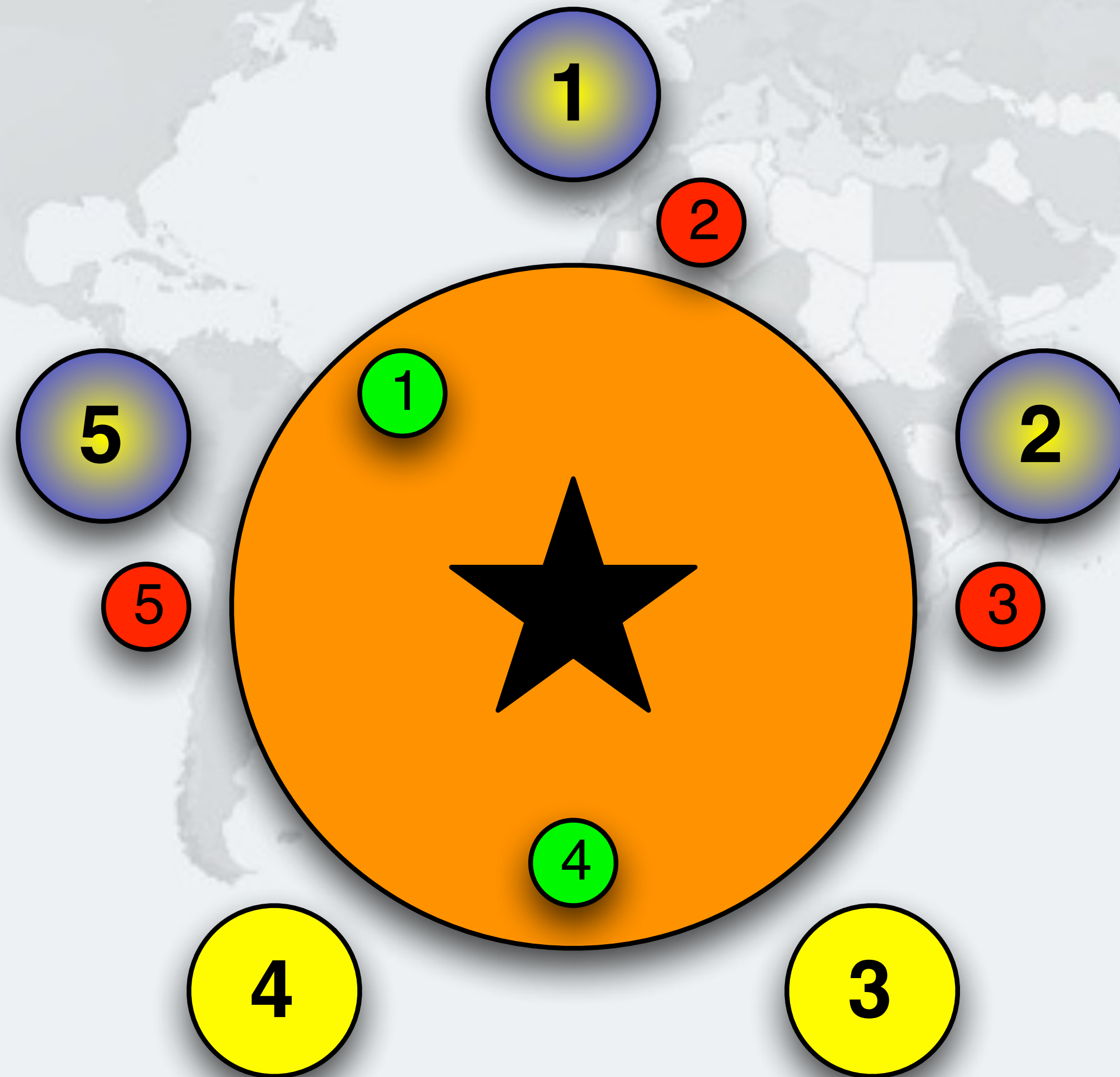


Philosopher 1 takes cup 2

- **Must take the cup with the higher number first**
 - In this case cup 2



Philosopher 2 takes cup 3

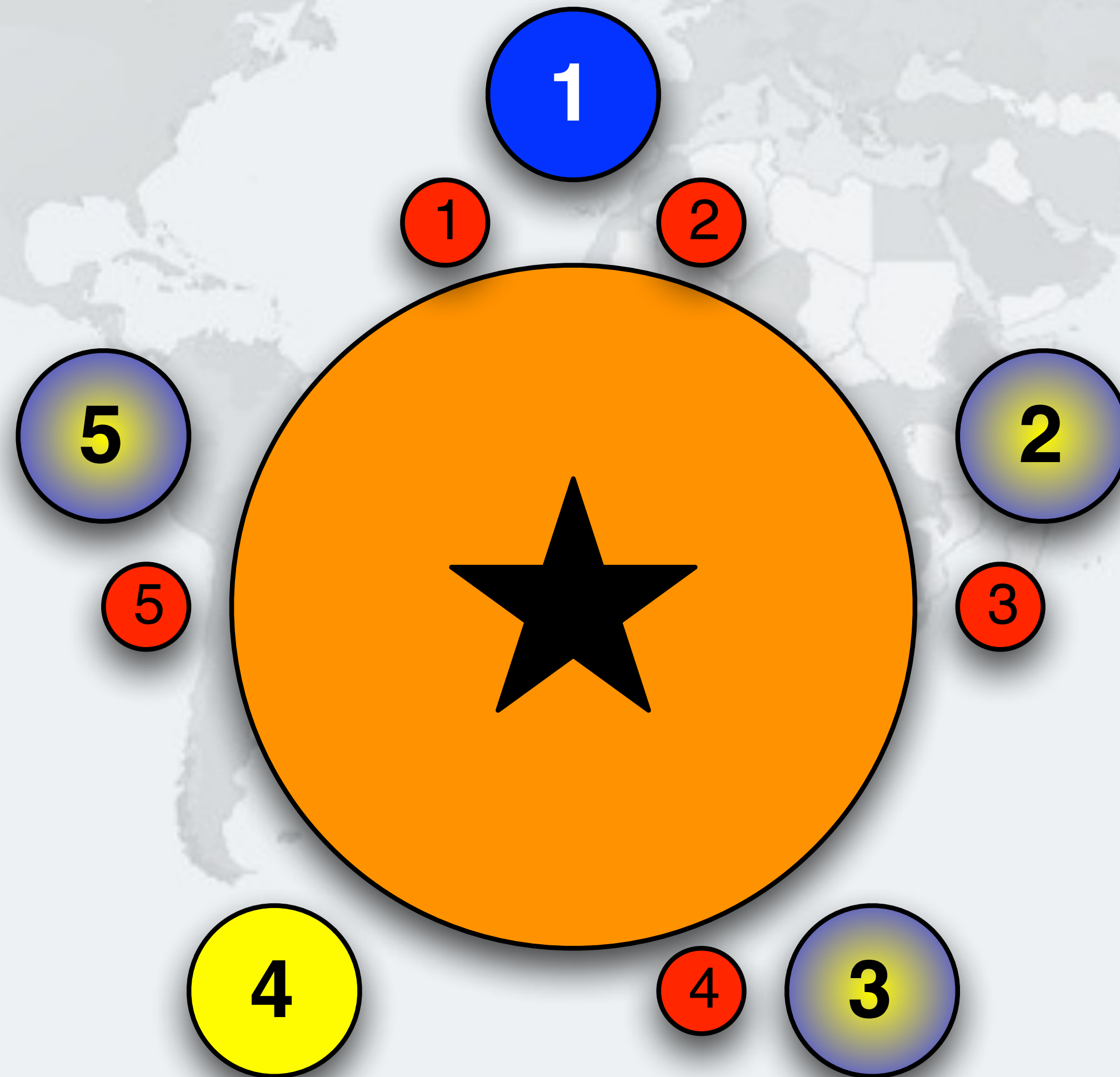


Philosopher 3 takes cup 4

- Note that philosopher 4 is prevented from holding one cup

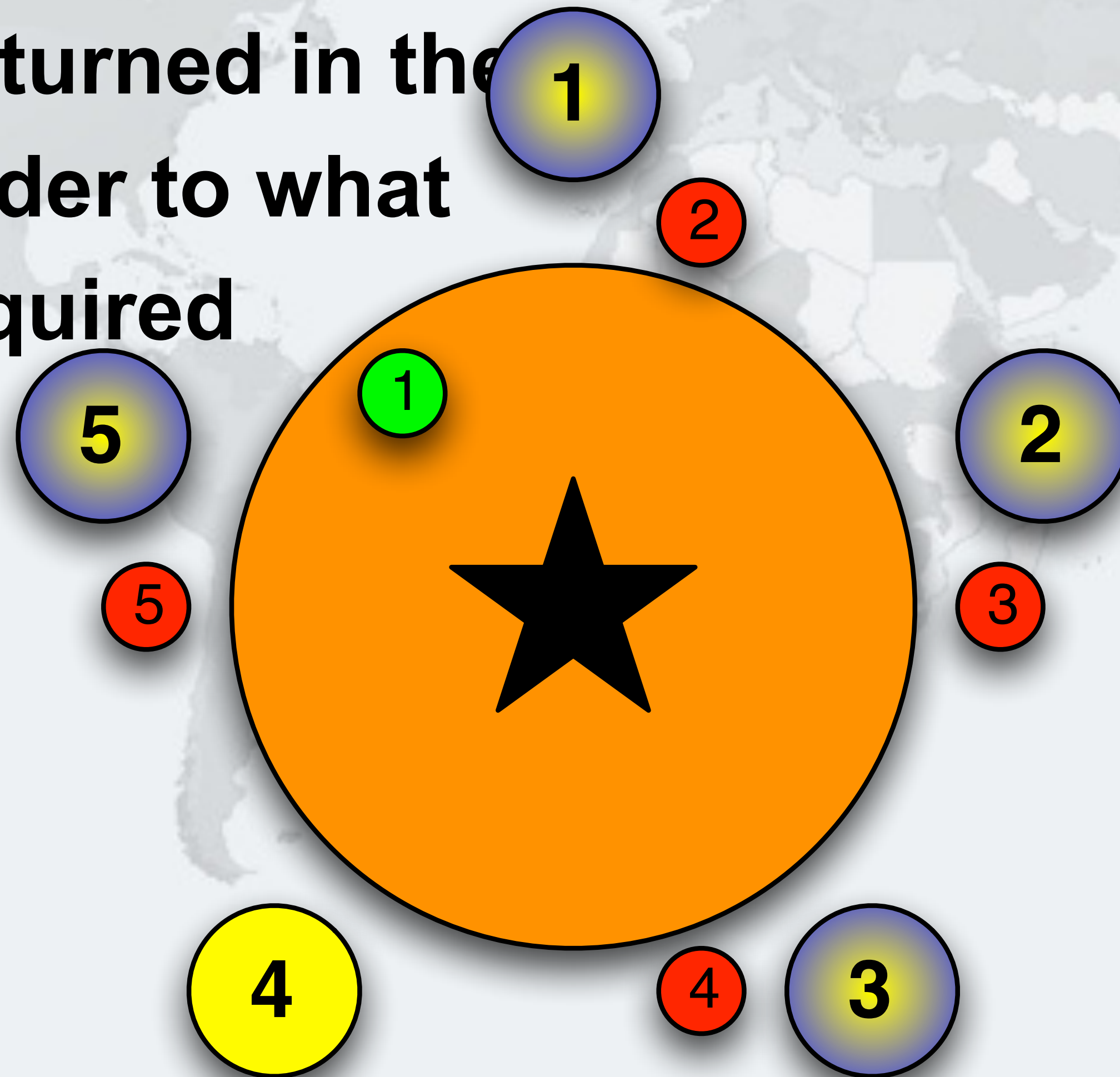


Philosopher 1 takes cup 1 - Drinking

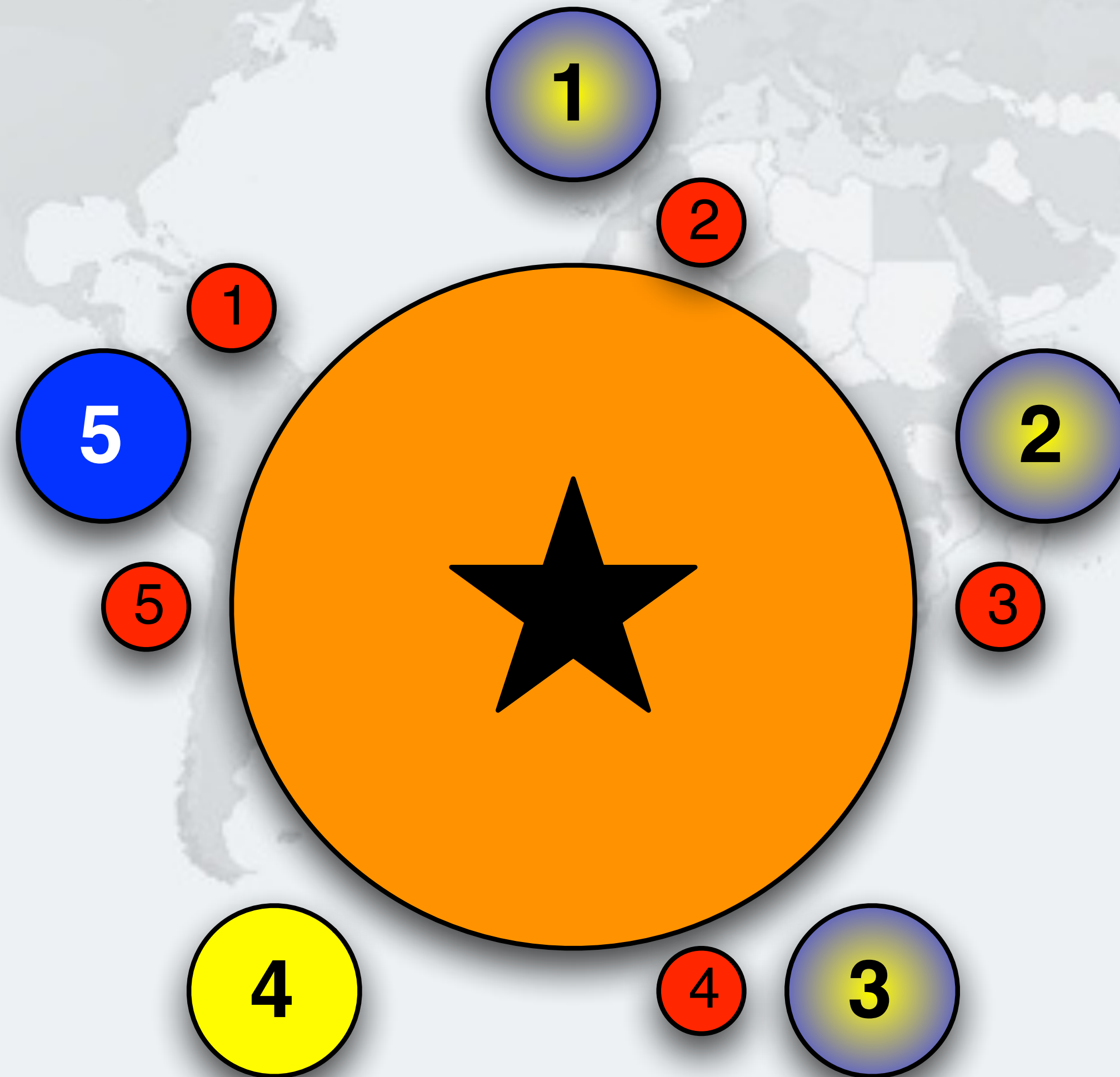


Philosopher 1 returns Cup 1

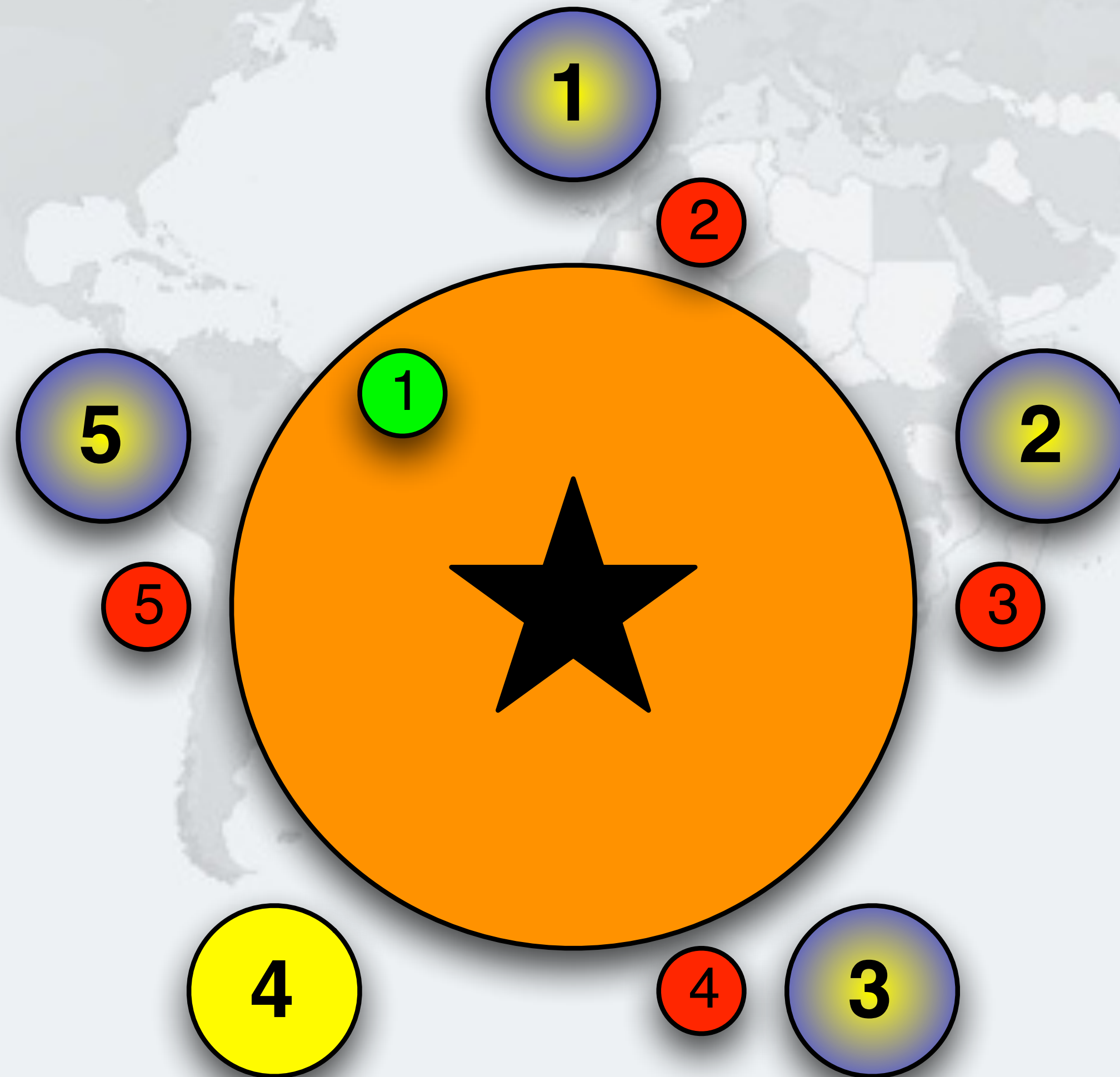
- **Cups are returned in the opposite order to what they are acquired**



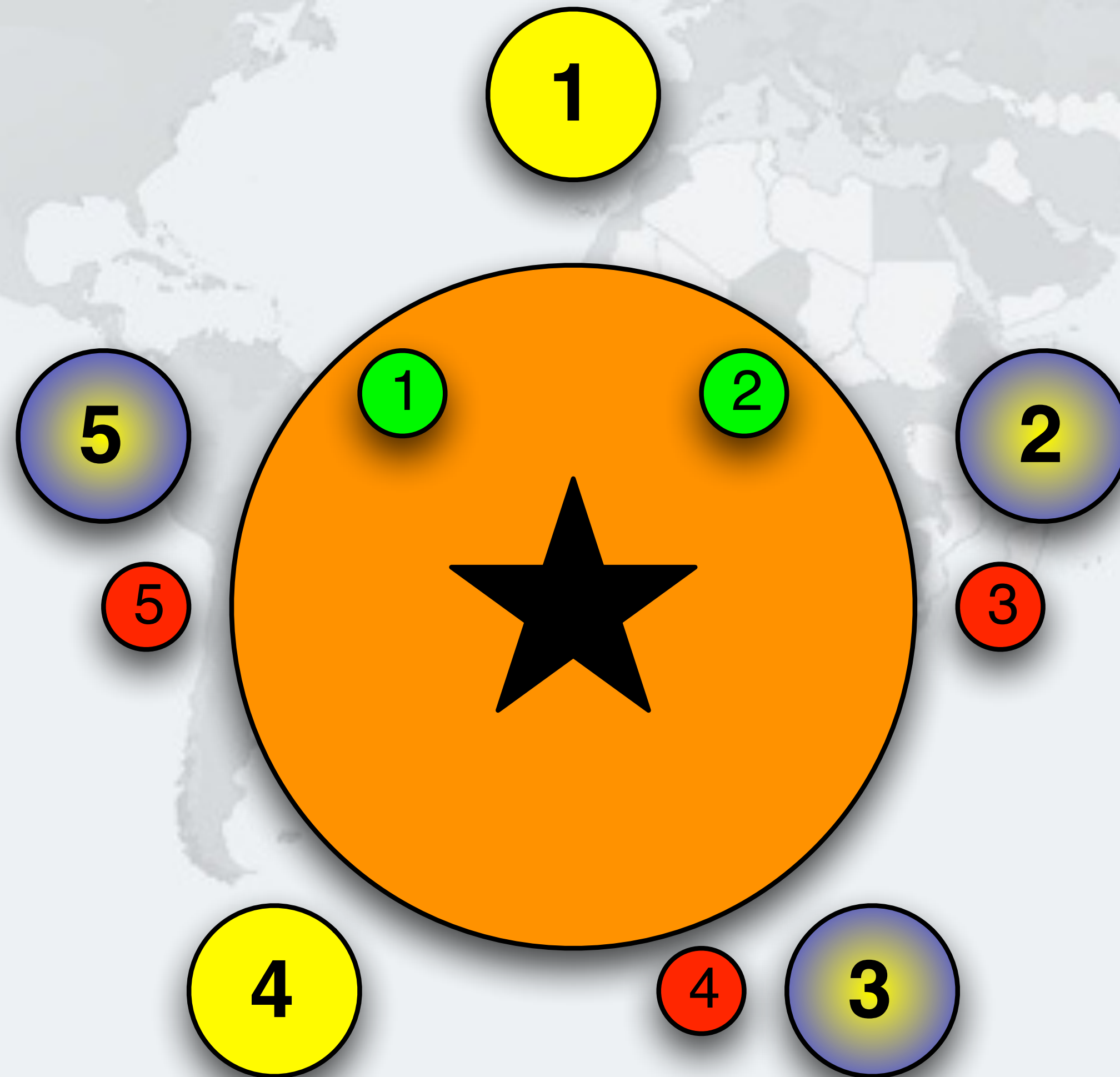
Philosopher 5 takes cup 1 - Drinking



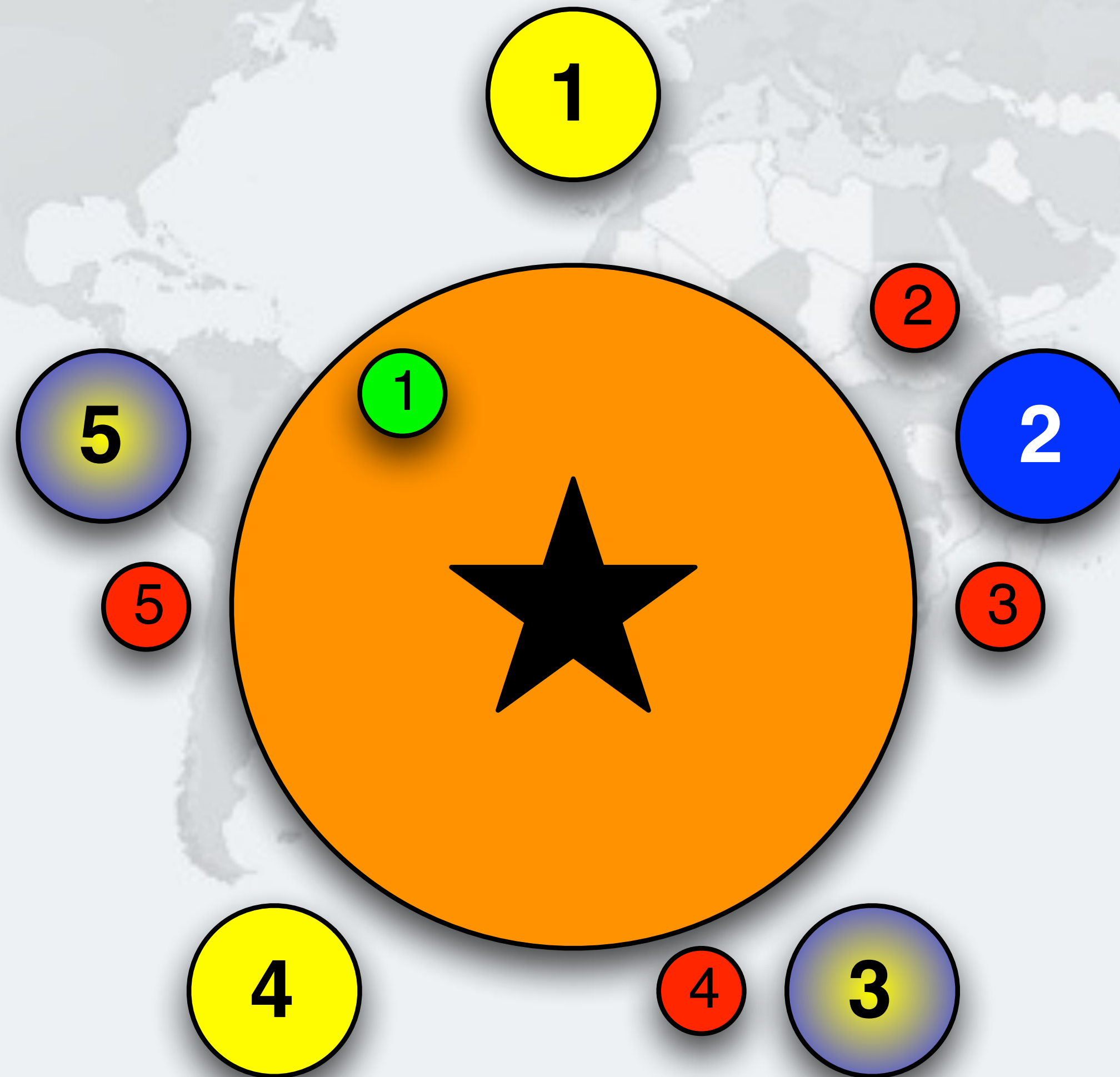
Philosopher 5 returns cup 1



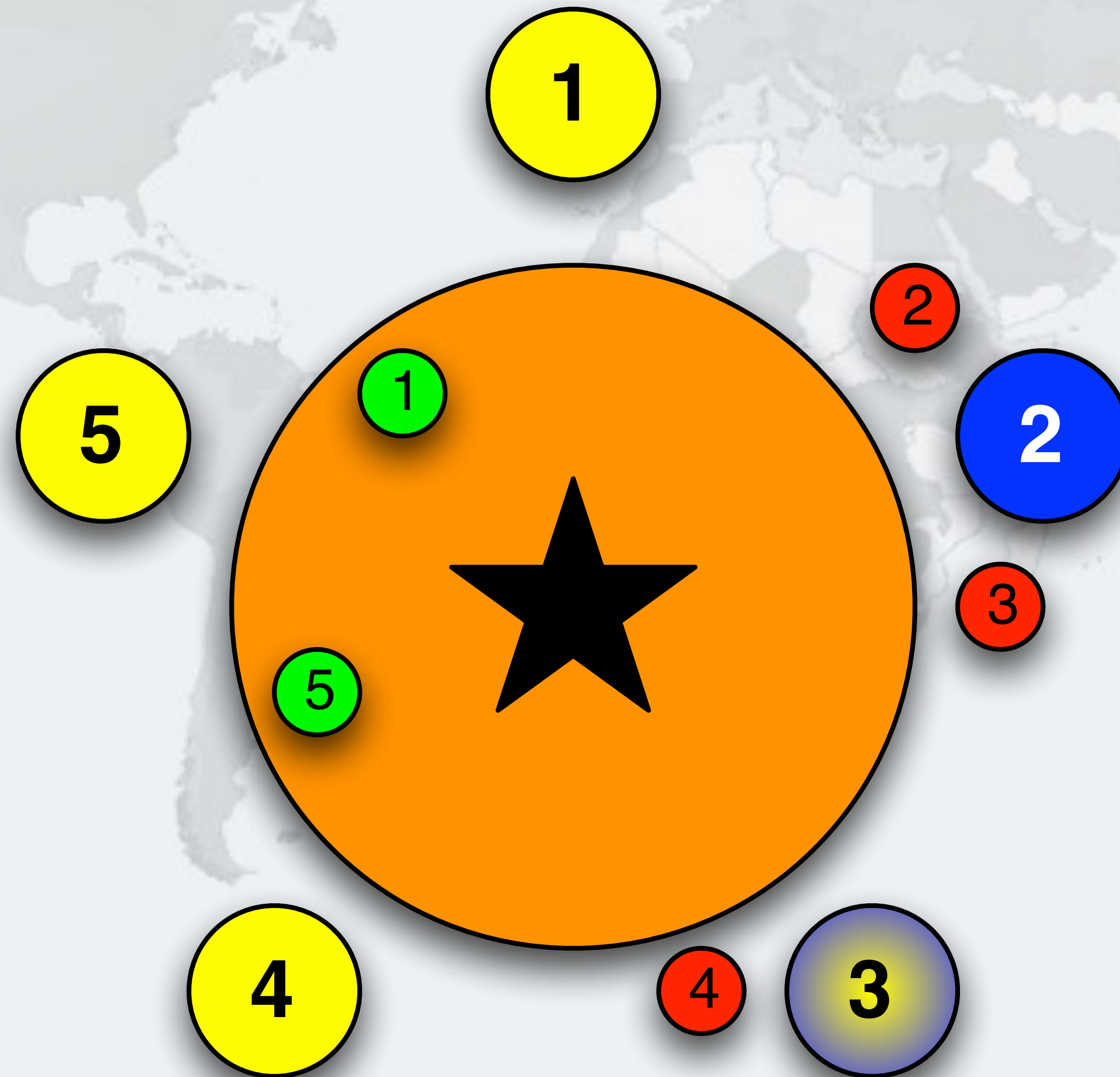
Philosopher 1 returns cup 2



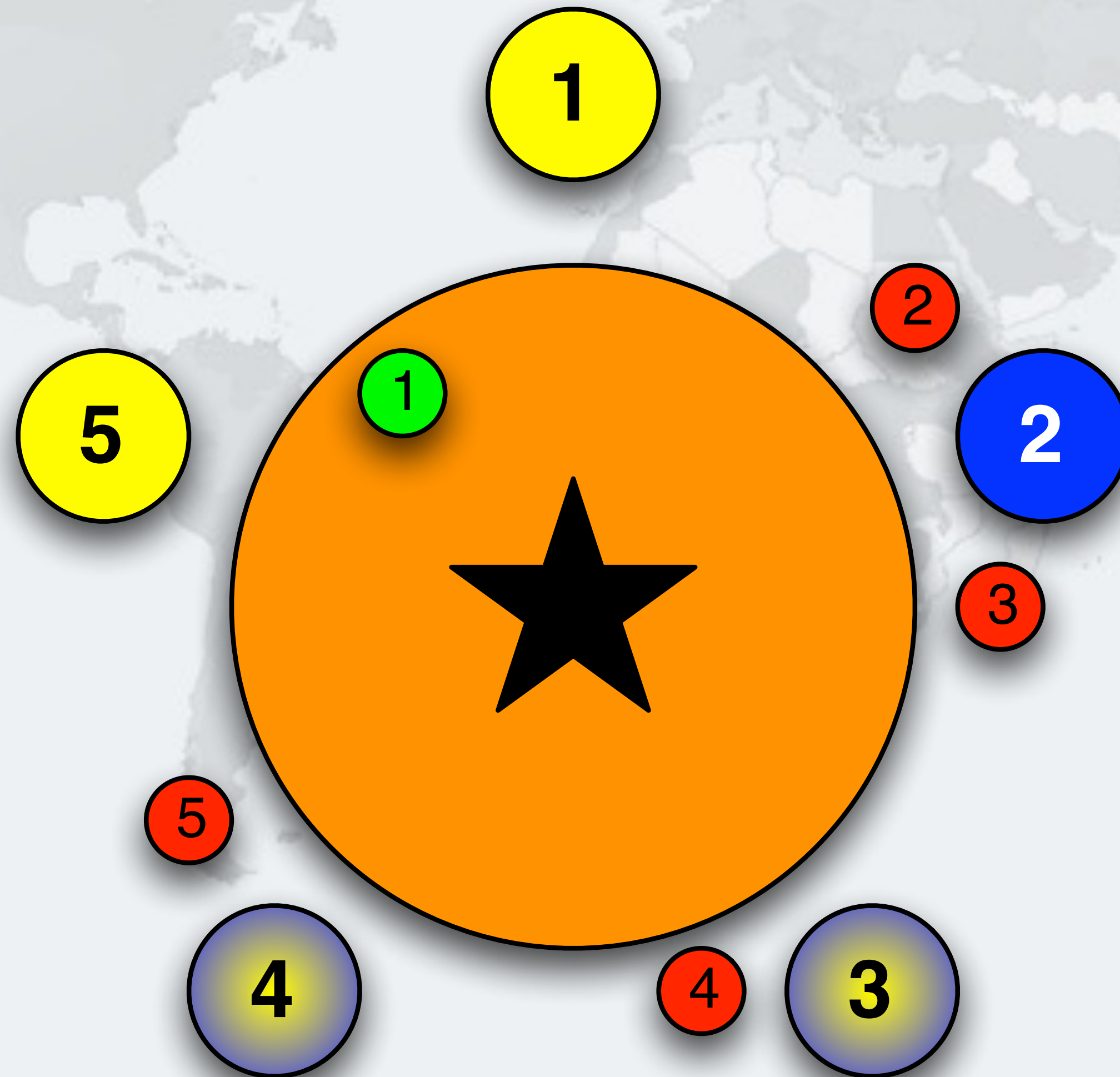
Philosopher 2 takes cup 2 - Drinking



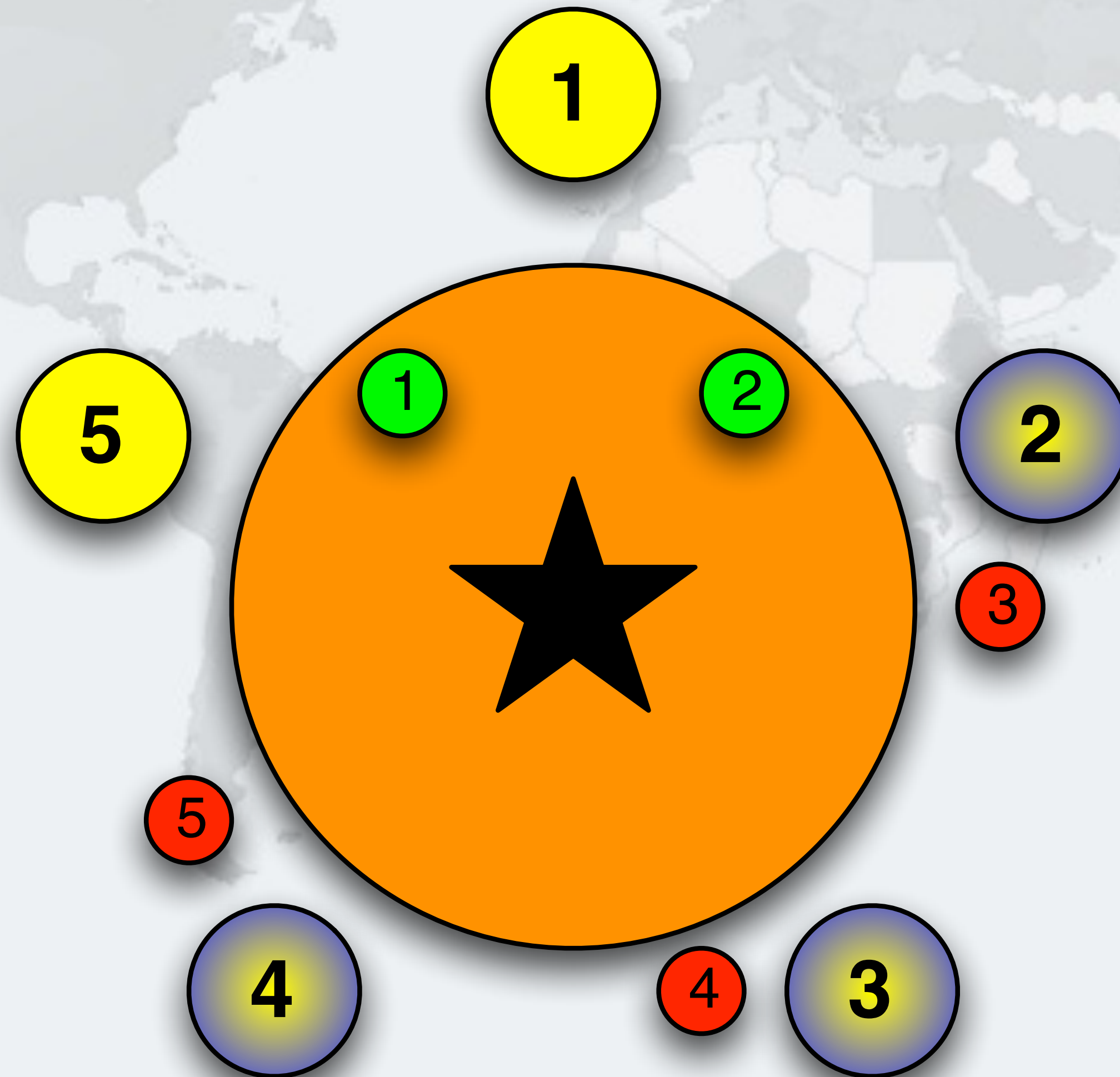
Philosopher 5 returns cup 5



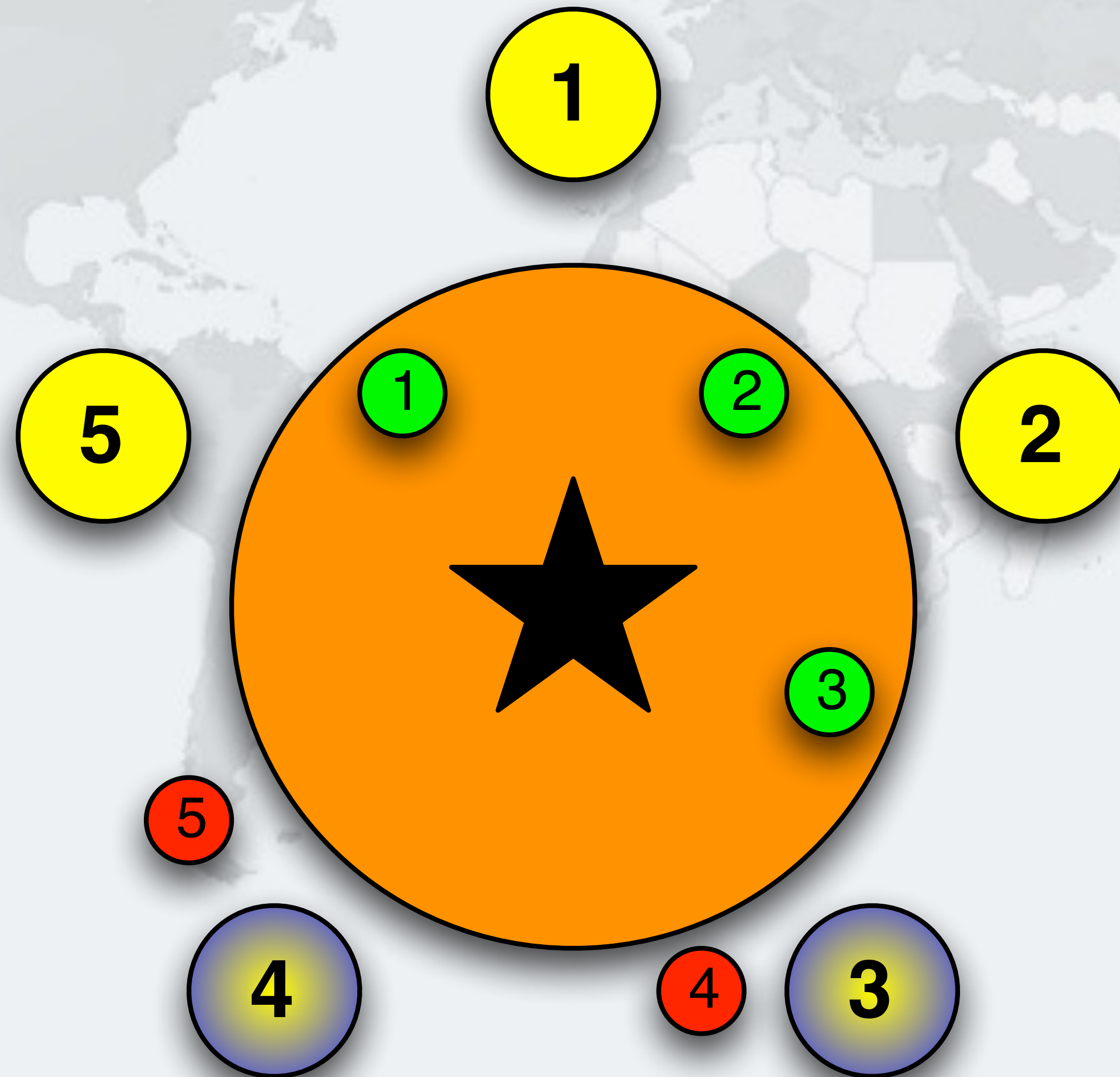
Philosopher 4 takes cup 5



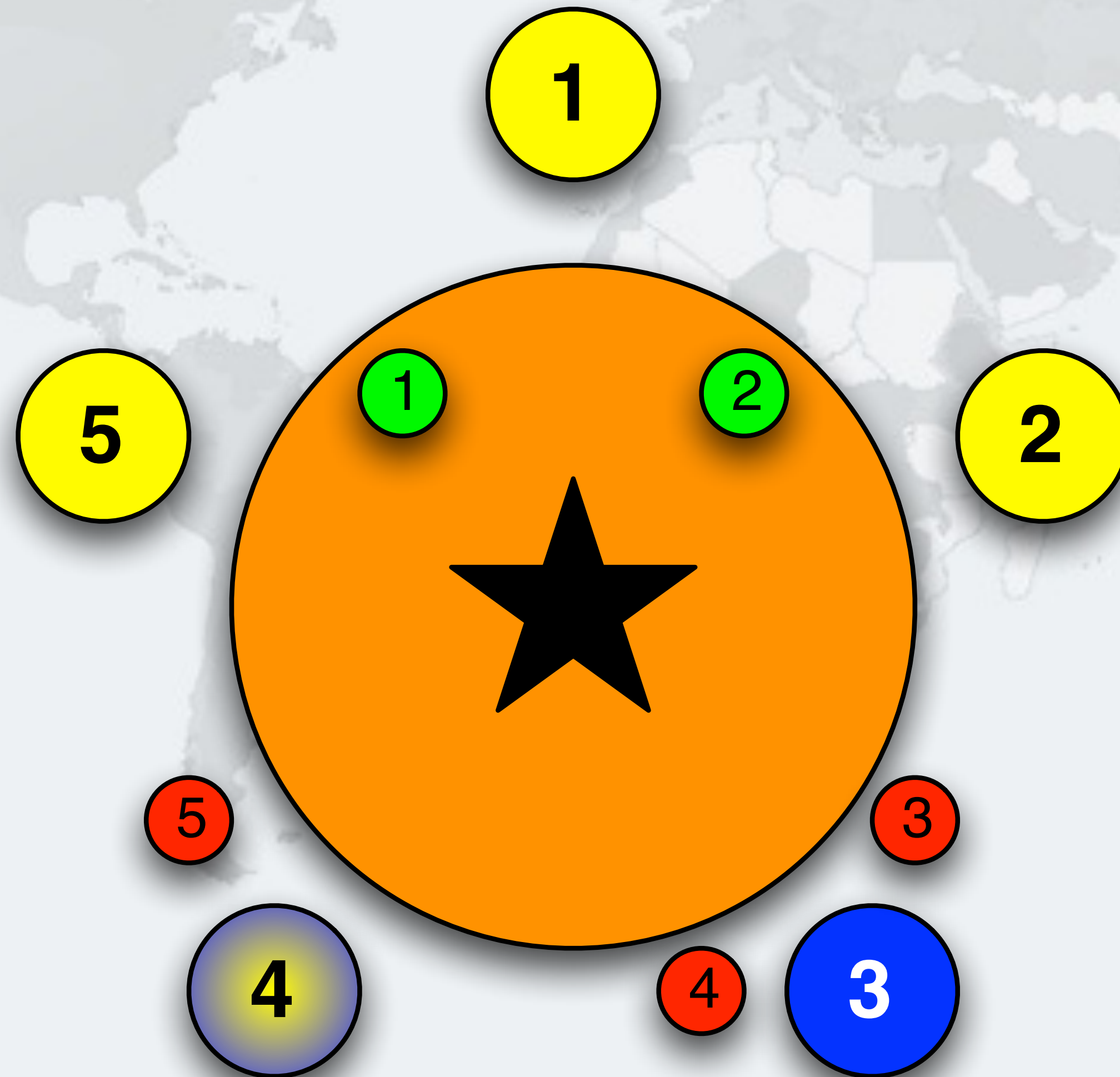
Philosopher 2 returns cup 2



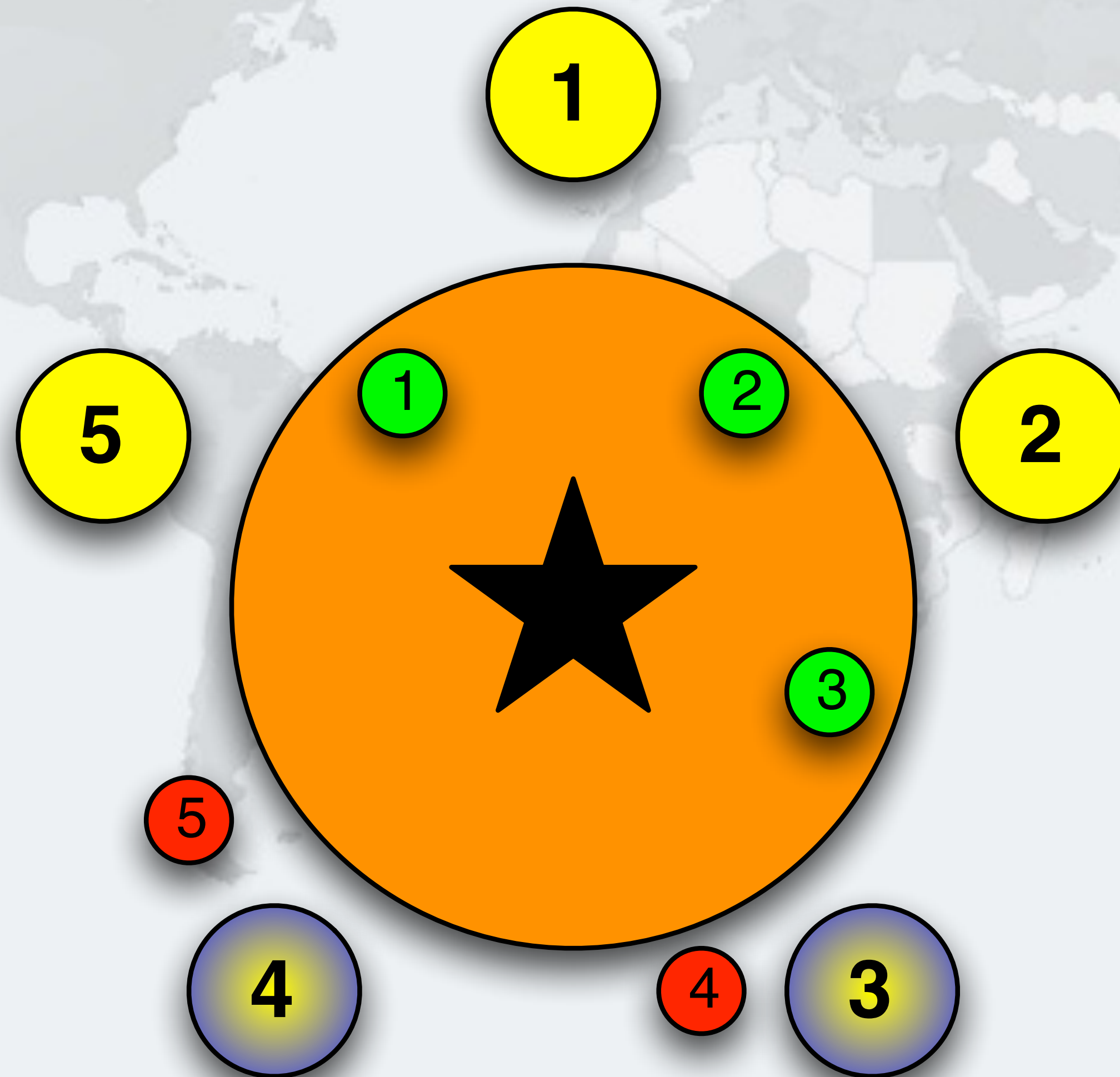
Philosopher 2 returns cup 3



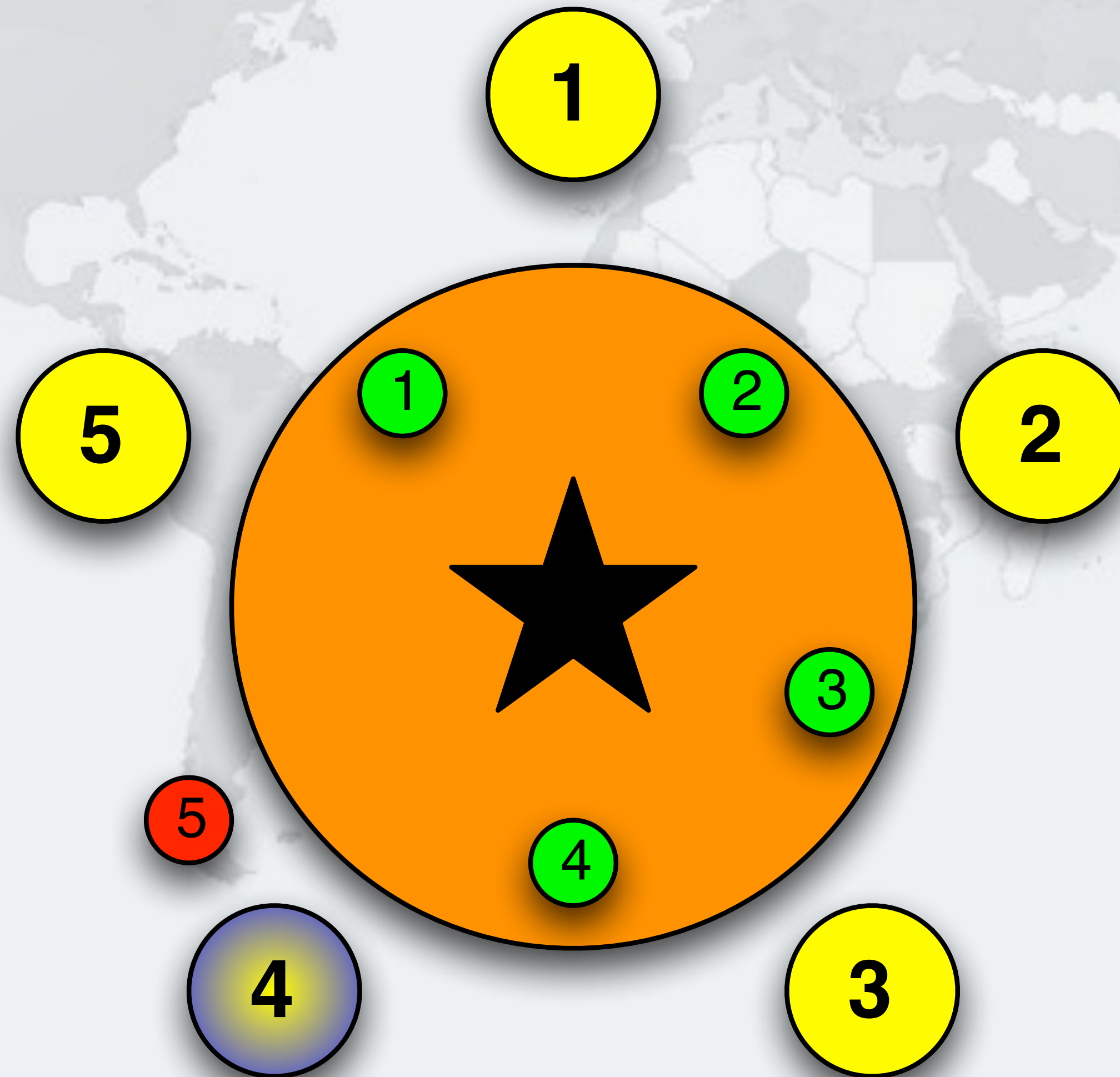
Philosopher 3 takes cup 3 - Drinking



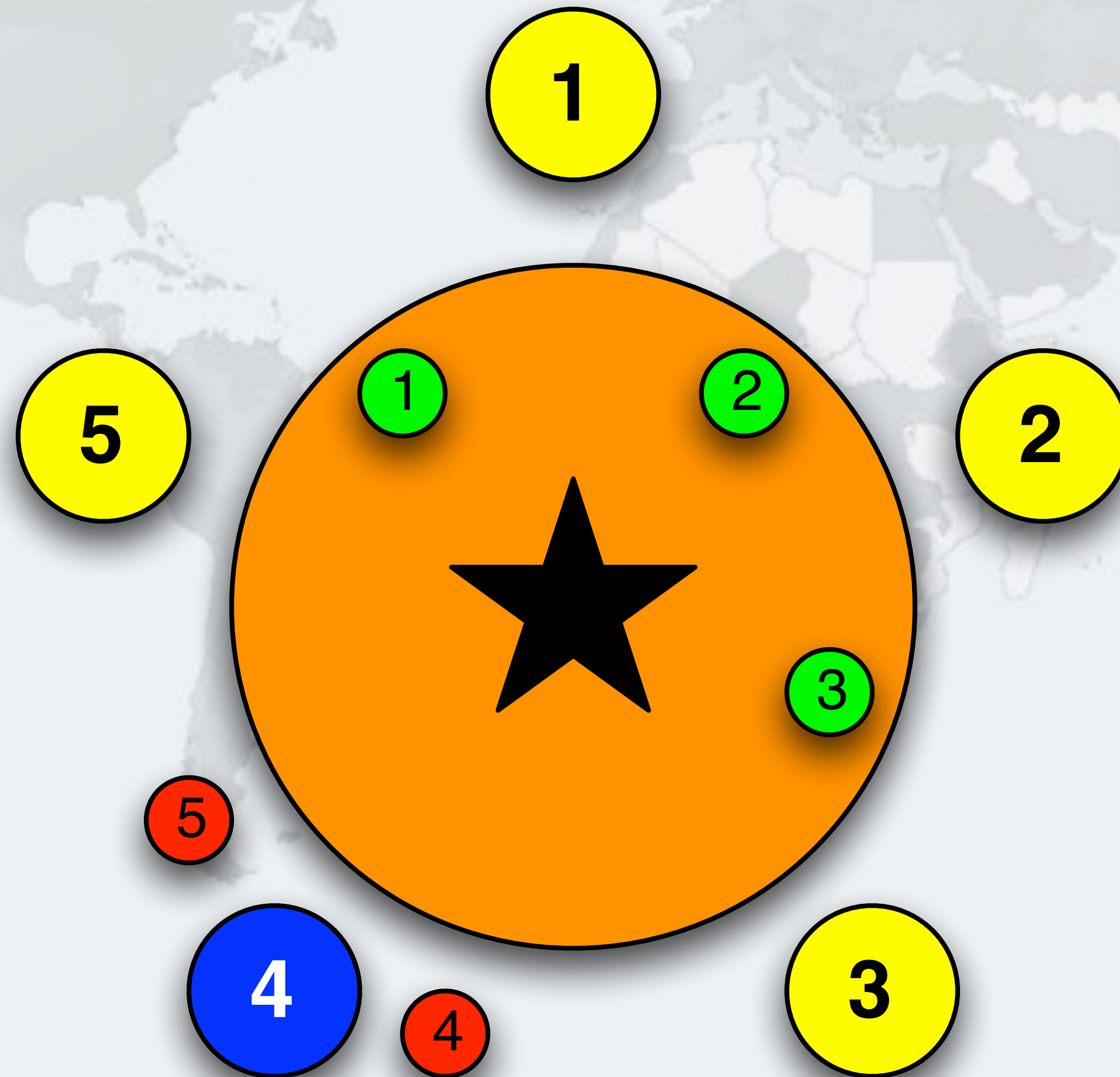
Philosopher 3 Returns cup 3



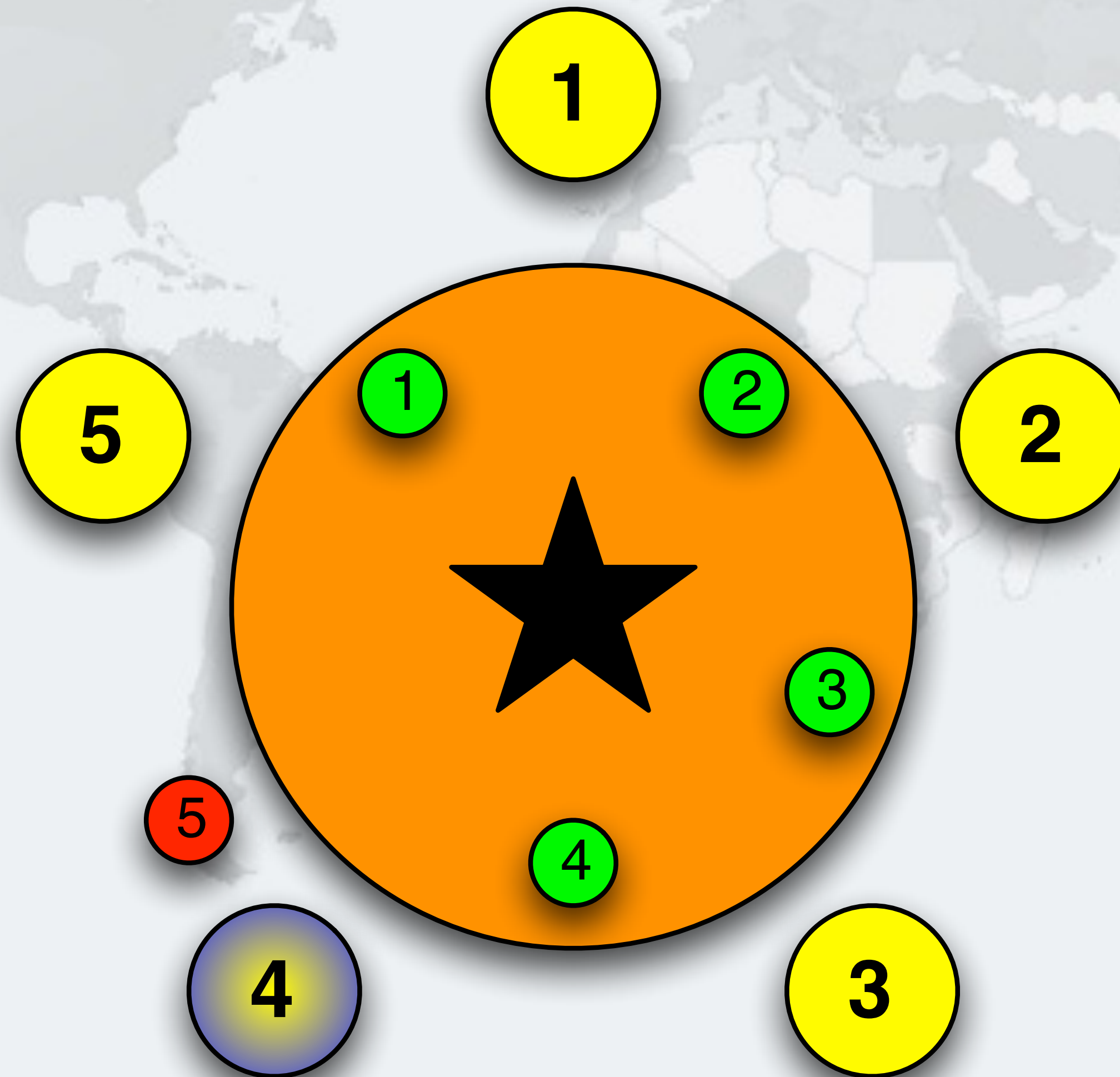
Philosopher 3 Returns cup 4



Philosopher 4 takes cup 4 - Drinking

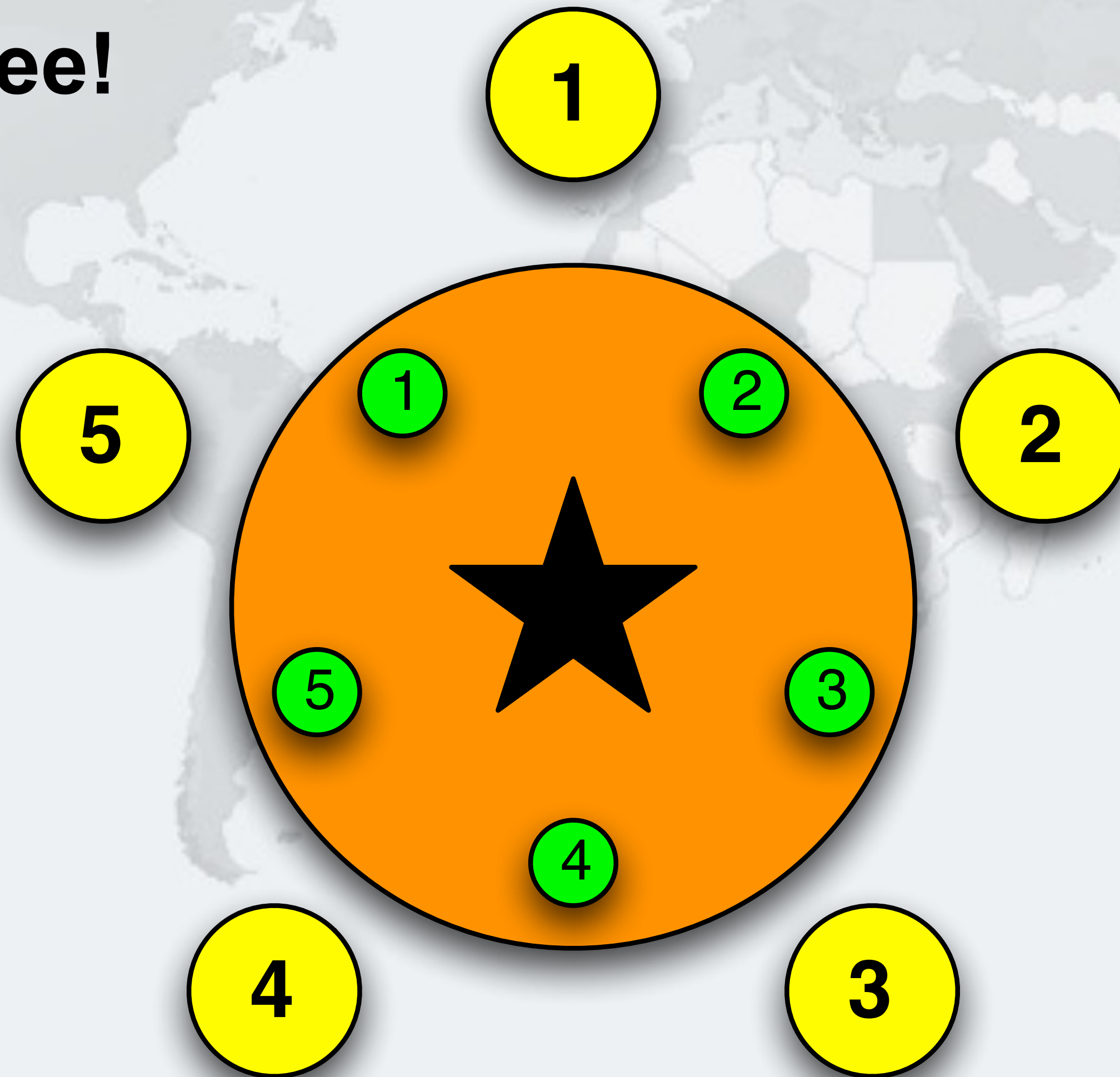


Philosopher 4 Returns cup 4



Philosopher 4 Returns cup 5

- **Deadlock free!**



Deadlock is avoided

- **Impossible for all philosophers to hold one cup**

Tools jstack and jps

- **We have command line tools**
 - **jps**
 - **shows your Java process ids**
 - **jstack -l pid**
 - **shows what your JVM is currently doing**
 - **Tools are in your jdk/bin directory**

Lab 1 Exercise



Lab1 - Save our philosophers

- **Define a global order on the locks**
 1. **Run `eu.javaspecialists.deadlock.lab1.Main`**
 2. **Make the `Krasi` object implement `Comparable`**
 3. **Lock first on bigger, than on smaller `Krasi`**
 4. **Verify that the deadlock has now disappeared**

Lab 2: Deadlock resolution by tryLock

Avoiding Liveness Hazards



Lab 2: Deadlock resolution by tryLock

- **Same problem as in Lab 1**
 - **But our solution will be different**
- **Instead of a global order on the locks**
 - **We lock the first lock**
 - **We then try to lock the second lock**
 - **If we can lock it, we start drinking**
 - **If we cannot, we back out completely and try again**
 - **What about starvation or livelock?**

Lock and ReentrantLock

- **The Lock interface offers different ways of locking:**
 - **Unconditional, polled, timed and interruptible**

```
public interface Lock {  
    void lock();  
    void lockInterruptibly() throws InterruptedException;  
    boolean tryLock();  
    boolean tryLock(long timeout, TimeUnit unit)  
        throws InterruptedException;  
    void unlock();  
    Condition newCondition();  
}
```

- **Lock implementations must have same memory-visibility semantics as intrinsic locks (synchronized)**

ReentrantLock Implementation

- **Like synchronized, it offers reentrant locking semantics**
- **Also, we can interrupt threads that are waiting for locks**
 - **Actually, the ReentrantLock never causes the thread to be BLOCKED, but always WAITING**
 - **If we try to acquire a lock unconditionally, interrupting the thread will simply go back into the WAITING state**
 - **Once the lock has been granted, the thread interrupts itself**

Using the explicit lock

- **We have to call `unlock()` in a `finally` block**

- Every time, without exception

- There are FindBugs detectors that will look for forgotten

"unlocks" `private final Lock lock = new ReentrantLock();`

```
public void update() {  
    lock.lock(); // this should be before try  
    try {  
        // update object state  
        // catch exceptions and restore  
        // invariants if necessary  
    } finally {  
        lock.unlock();  
    }  
}
```


Polled lock acquisition

- **Instead of unconditional lock, we can tryLock()**

```
if (lock.tryLock()) {  
    try {  
        balance = balance + amount;  
    } finally {  
        lock.unlock();  
    }  
} else {  
    // alternative path  
}
```

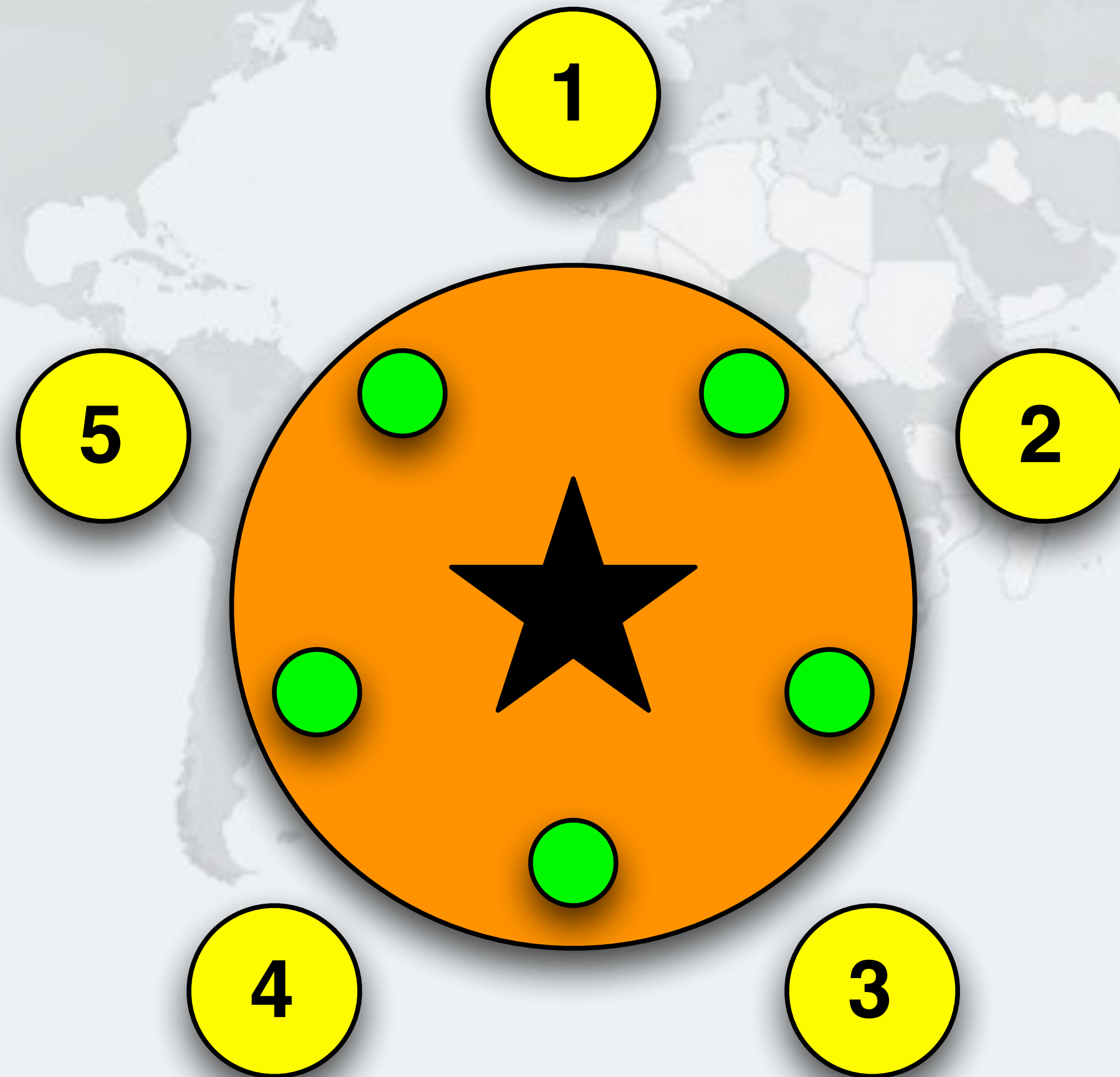
Using try-Lock to avoid deadlocks

- **Deadlocks happen when we lock multiple locks in different orders**
- **We can avoid this by using tryLock()**
 - If we do not get lock, sleep for a random time and then try again
 - Must release *all* held locks, or our deadlocks become livelocks
- **This is possible with synchronized, see my newsletter**
 - <http://www.javaspecialists.eu/archive/Issue194.html>

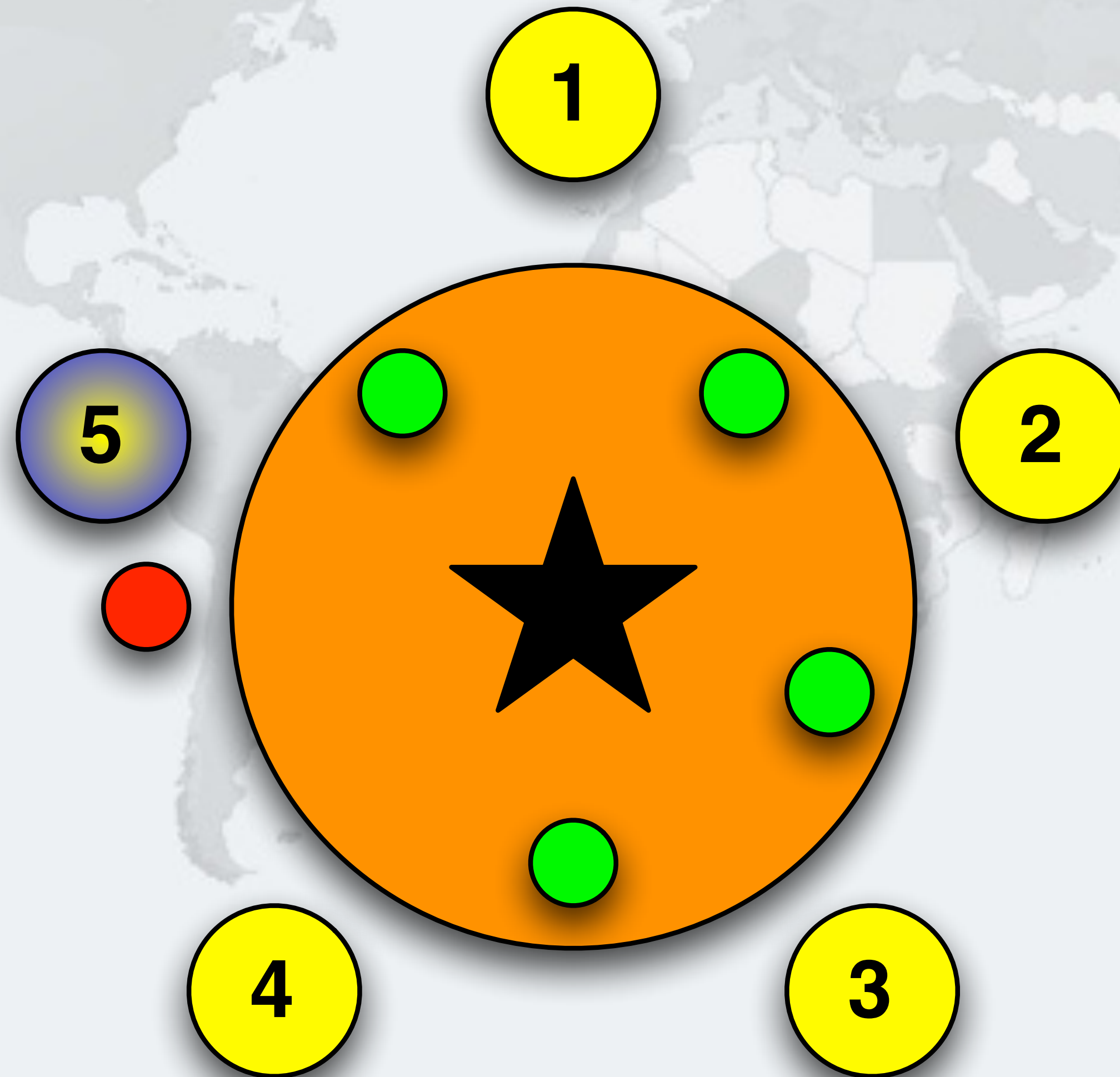
Using Trylock() To Avoid Deadlocks

```
public void drink() {  
    while (true) {  
        left.lock();  
        try {  
            if (right.tryLock()) {  
                try {  
                    // now we can finally drink and then return  
                    return;  
                } finally {  
                    right.unlock();  
                }  
            }  
        } finally {  
            left.unlock();  
        }  
    }  
}
```

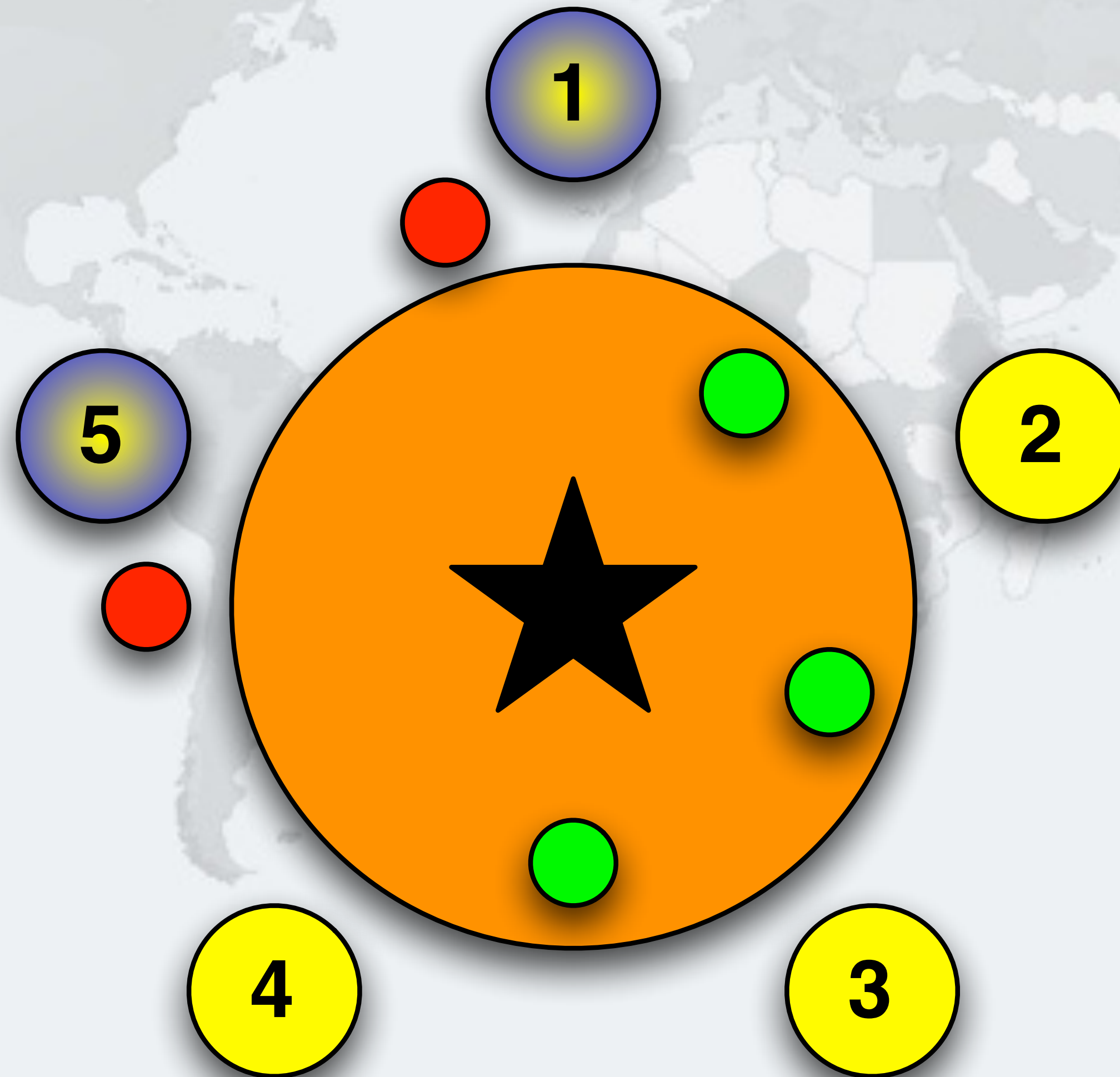
Deadlock is prevented in this design



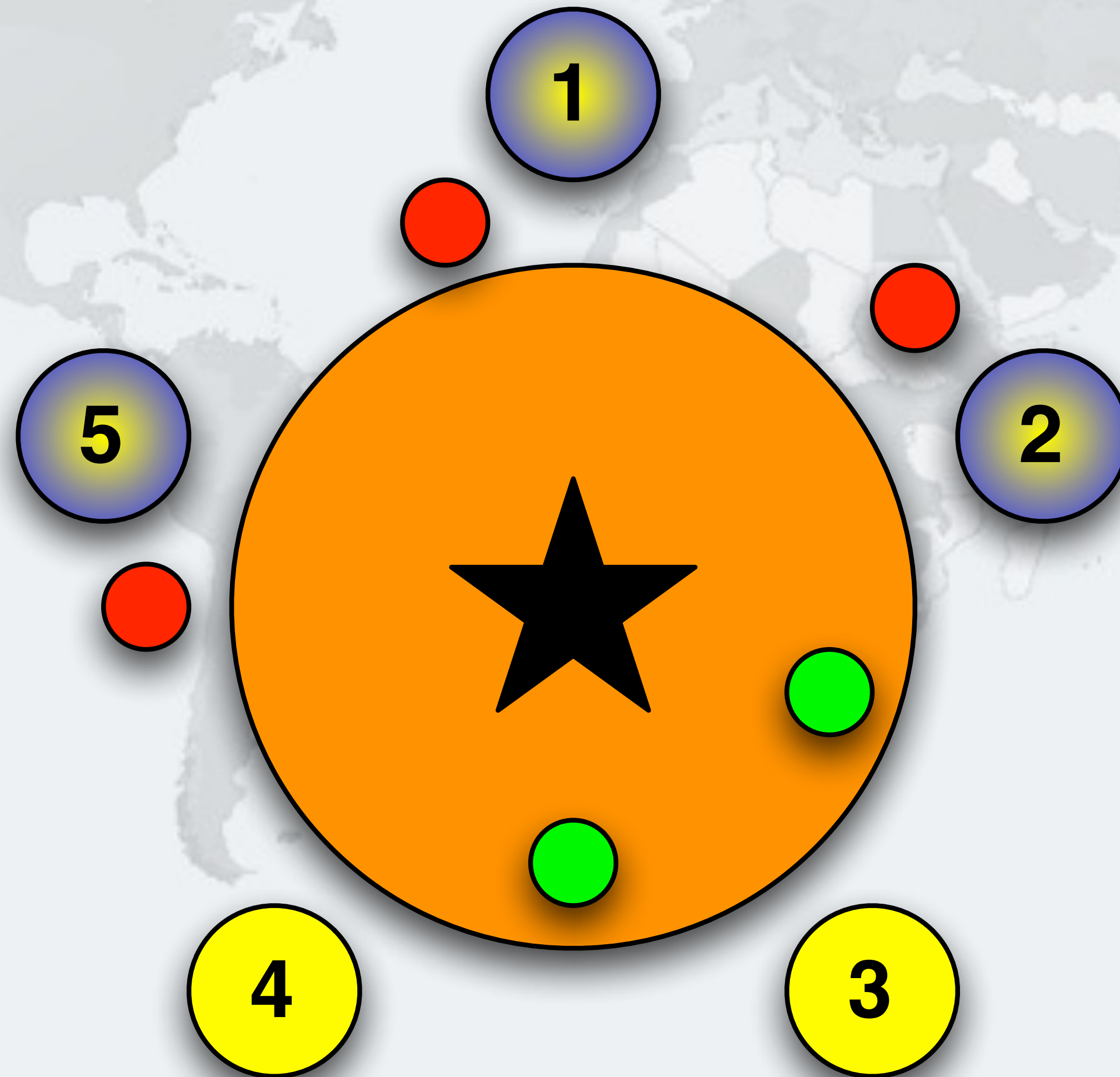
Philosopher 5 wants to drink, takes right cup



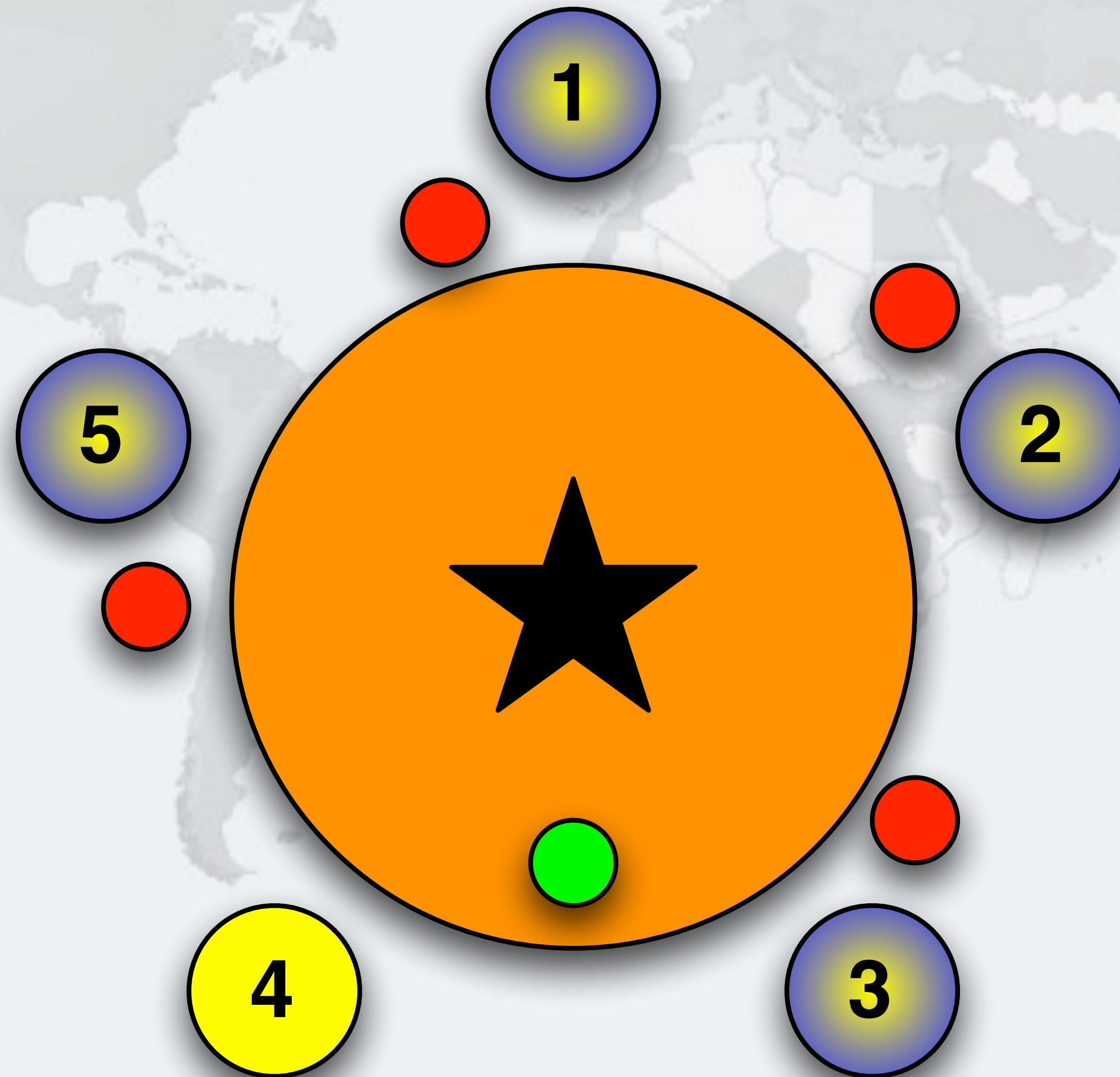
Philosopher 1 wants to drink, takes right cup



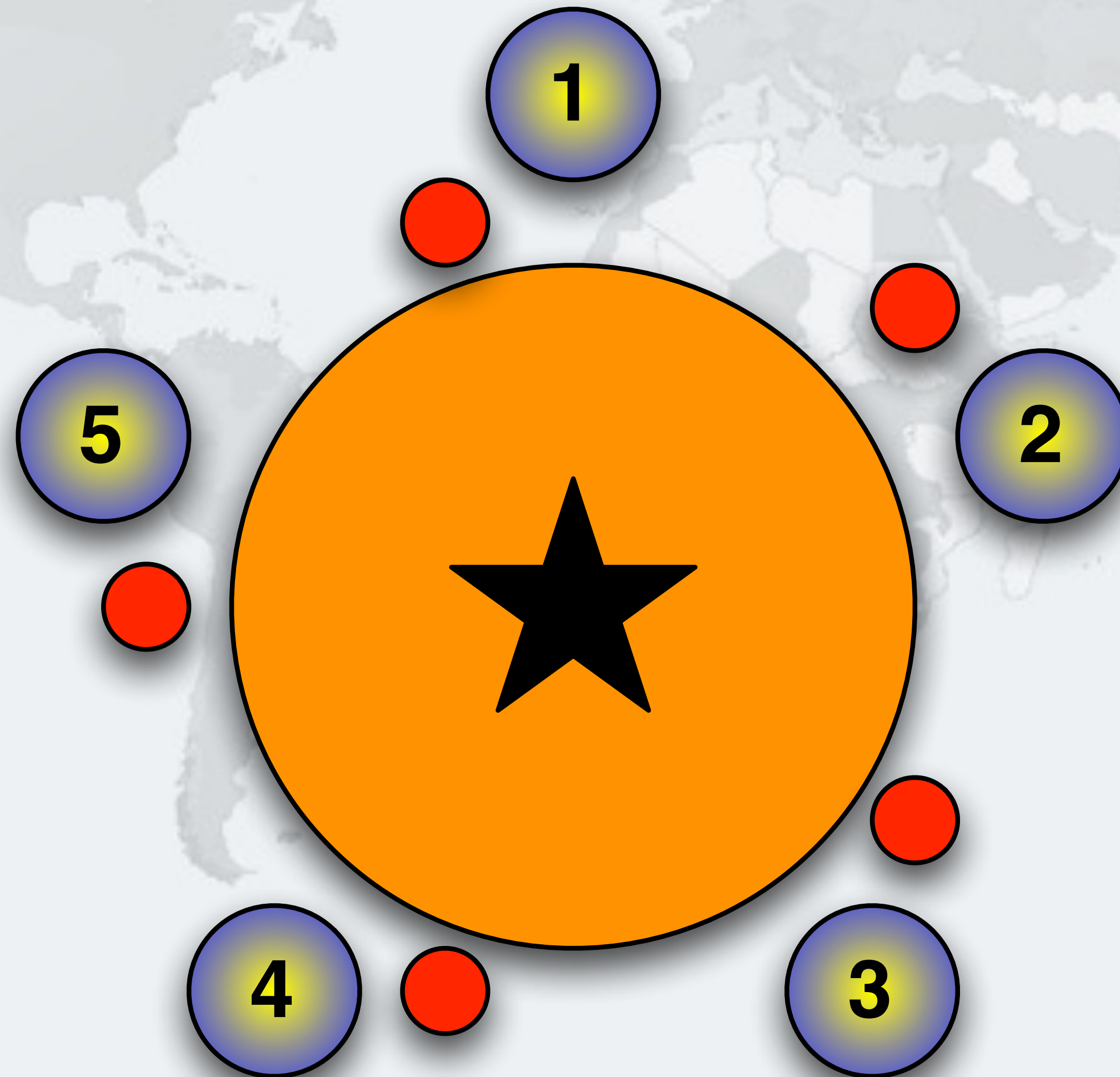
Philosopher 2 wants to drink, takes right cup



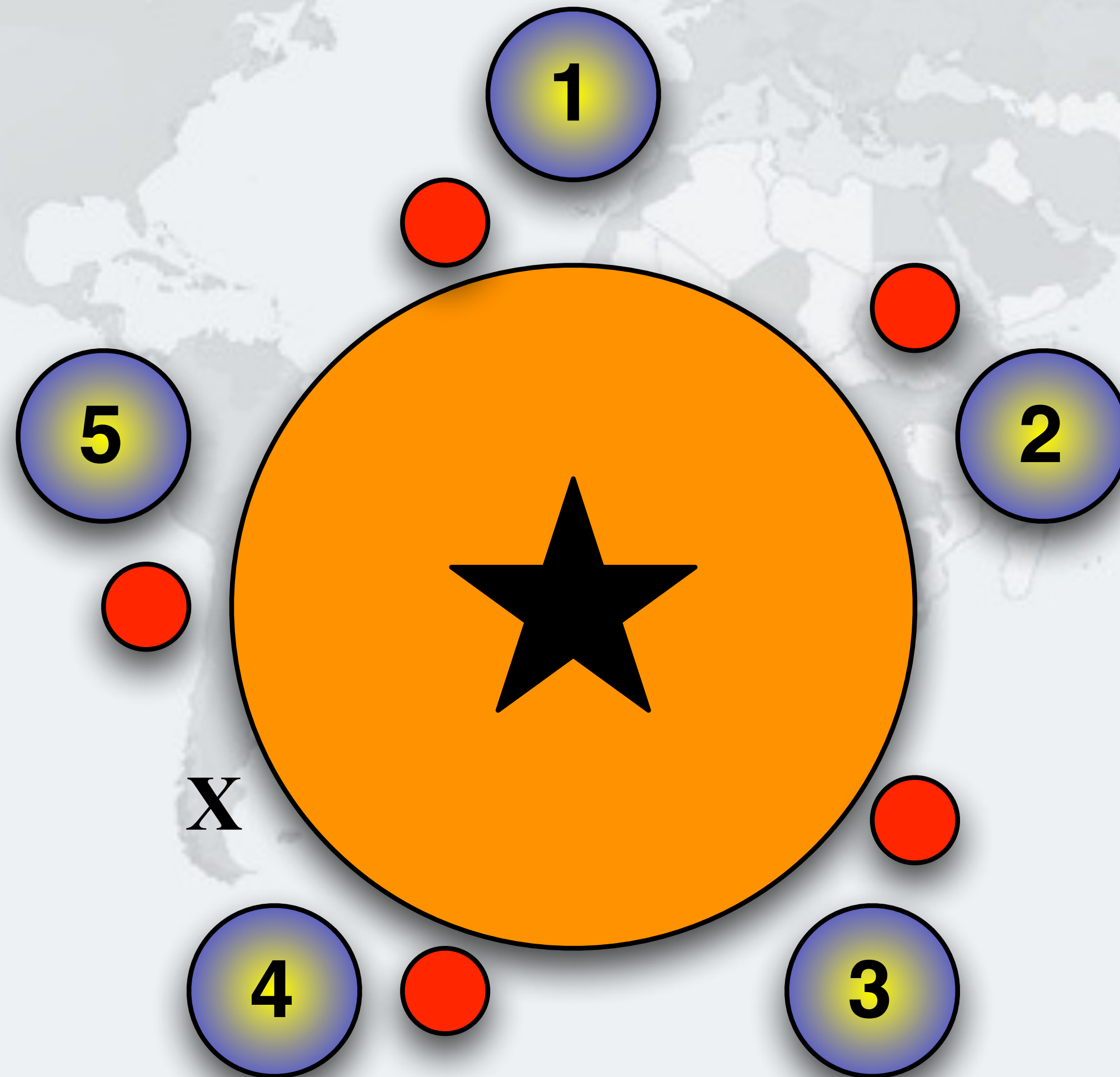
Philosopher 3 wants to drink, takes right cup



Philosopher 4 wants to drink, takes right cup

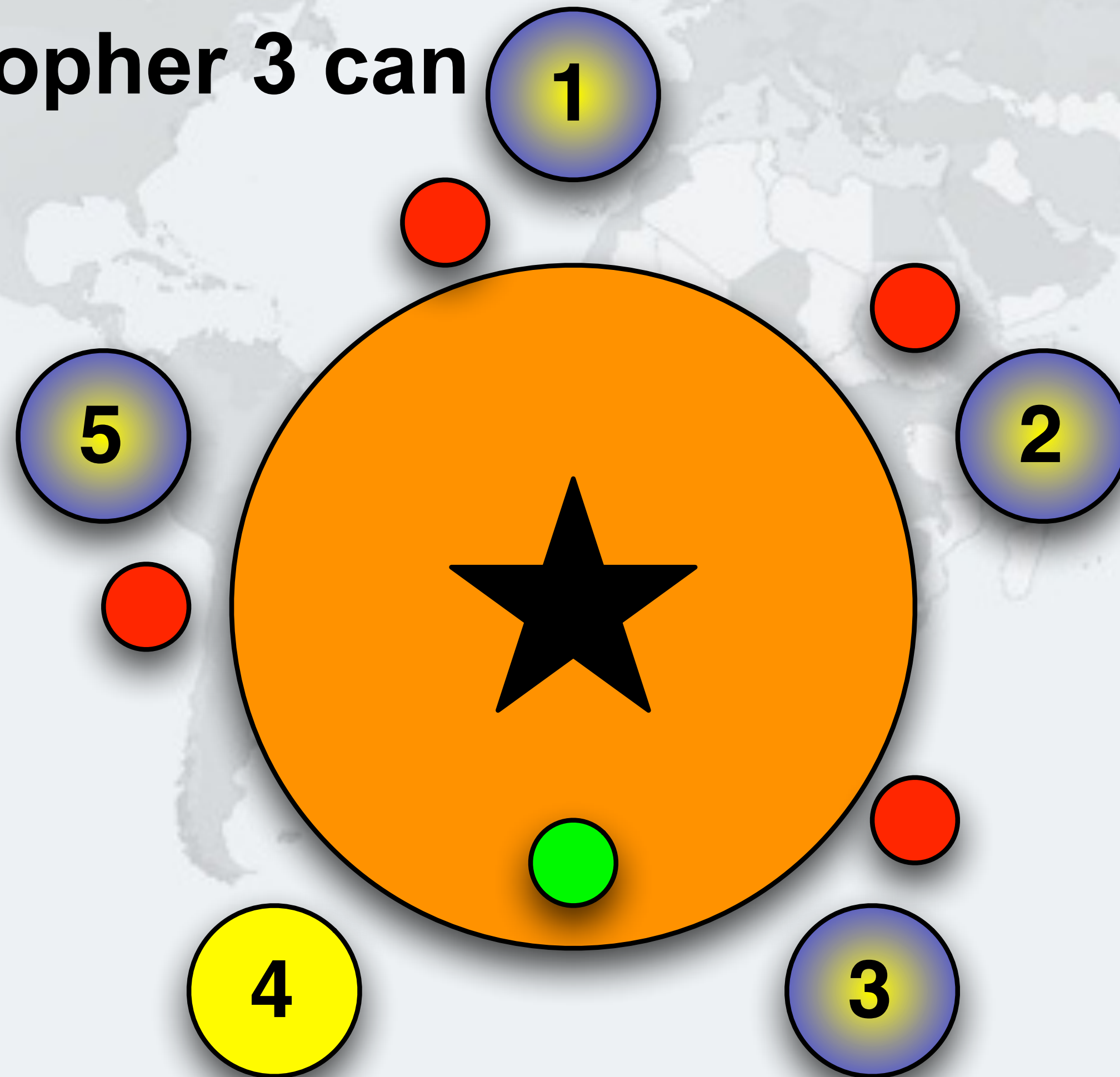


Philosopher 4 tries to lock left, not available



Philosopher 4 Unlocks right again

- **Now Philosopher 3 can drink**



Lab 2 Exercise

Deadlock resolution with `tryLock()`



Lab2: Solving Deadlock with tryLock()

- **Run Main class to trigger deadlock**
 - **Capture a stack trace with `jstack -l pid`**
 - **Use `Lock.tryLock()` to avoid blocking on the inner lock**
 - **lock the right**
 - **tryLock the left**
 - **if success, then drink and unlock both**
 - **otherwise, unlock right only and retry**
 - **Verify that the deadlock has now disappeared**

Lab 3: Resource Deadlock

Avoiding Liveness Hazards



Lab 3: Resource Deadlock

- **Problem: threads are blocked waiting for a finite resource that never becomes available**
- **Examples:**
 - **Resources not being released after use**
 - **Running out of threads**
 - **Java Semaphores not being released**
 - **JDBC transactions getting stuck**
 - **Bounded queues or thread pools getting jammed up**

Challenge

- **Does not show up as a Java thread deadlock**
- **Problem thread could be in any state: RUNNABLE, WAITING, BLOCKED, TIMED_WAITING**

How to solve resource deadlocks

- **If you can reproduce the resource deadlock**
 - Take a thread dump shortly before the deadlock
 - Take another dump after the deadlock
 - Compare the two dumps
- **If you are already deadlocked**
 - Take a few thread dumps
 - Look for threads that don't move, but should

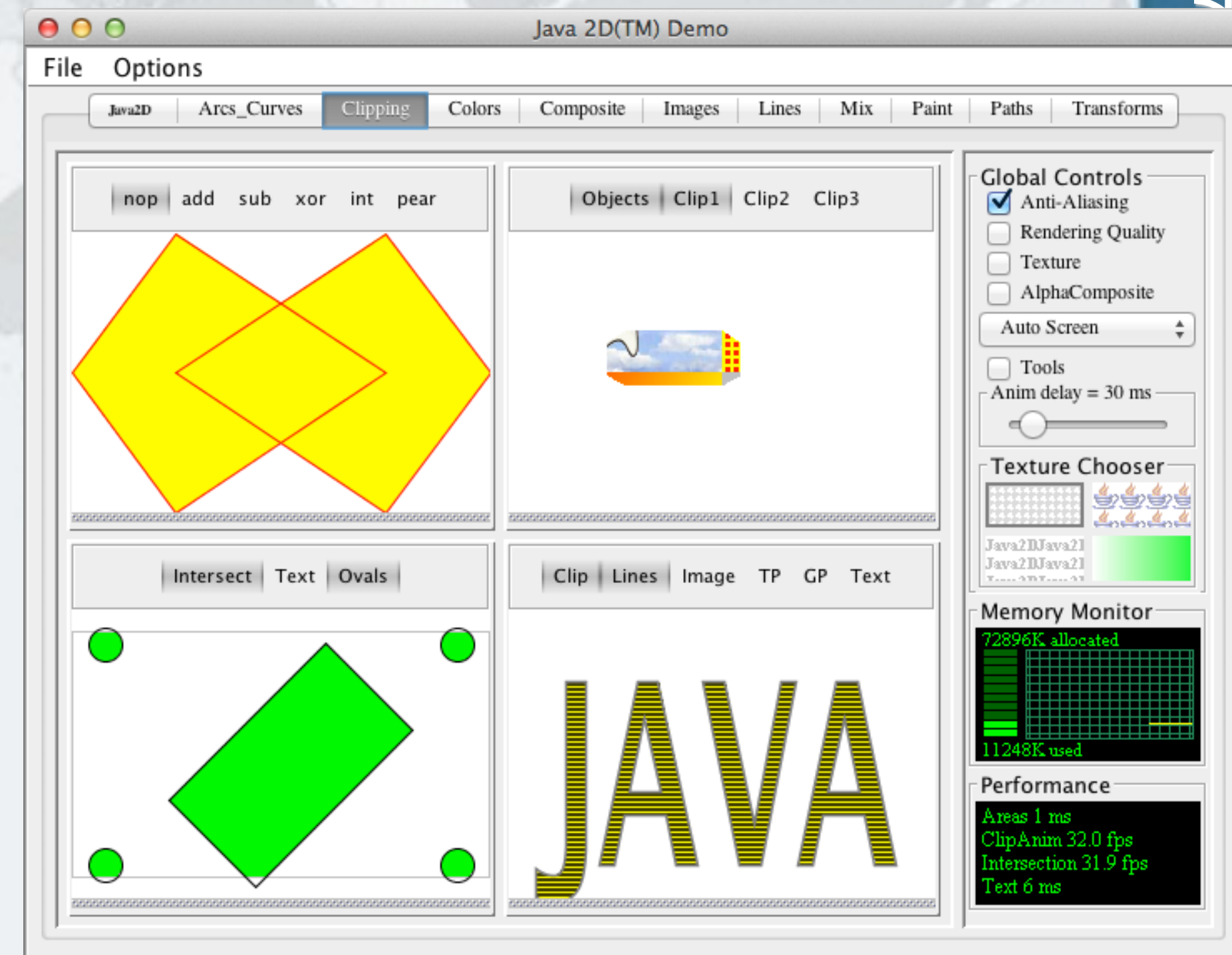
Lab 3 Exercise

Resource Deadlock



Lab3 Resource Deadlock

- **Start our modified Java2Demo**
 - Dump threads with `jstack -l`
 - Use Java2Demo for a while until it deadlocks
 - Get another thread dump and compare to the first one
 - This should show you where the problem is inside your code
 - Fix the problem and verify that it has been solved



Lab3 Exercise solution Explanation

- **Goal: Ensure that resources are released after use**
- **Diff between the two thread dumps using jps and jstack**

```
< at java.util.concurrent.locks.AbstractQueuedSynchronizer$ConditionObject.await(AbstractQueuedSynchronizer.java:2043)
< at java.awt.EventQueue.getNextEvent(EventQueue.java:531)
< at java.awt.EventDispatchThread.pumpOneEventForFilters(EventDispatchThread.java:213)
...
> at java.util.concurrent.locks.AbstractQueuedSynchronizer.parkAndCheckInterrupt(AbstractQueuedSynchronizer.java:834)
> at java.util.concurrent.locks.AbstractQueuedSynchronizer.doAcquireSharedInterruptibly(AbstractQueuedSynchronizer.java:994)
> at java.util.concurrent.locks.AbstractQueuedSynchronizer.acquireSharedInterruptibly(AbstractQueuedSynchronizer.java:1303)
> at java.util.concurrent.Semaphore.acquire(Semaphore.java:317)
> at eu.javaspecialists.deadlock.lab3.java2d.MemoryManager.gc(MemoryManager.java:56)
> at eu.javaspecialists.deadlock.lab3.java2d.MemoryMonitor$Surface.paint(MemoryMonitor.java:153)
```

- **Fault is probably in our classes, rather than JDK**

What Is Wrong With This Code?

```
/**
 * Only allow a maximum of 30 threads to call System.gc() at a time.
 */
public class MemoryManager extends Semaphore {
    private static final int MAXIMUM_NUMBER_OF_CONCURRENT_GC_CALLS = 30;

    public MemoryManager() {
        super(MAXIMUM_NUMBER_OF_CONCURRENT_GC_CALLS);
    }

    public void gc() {
        try {
            acquire();
            try {
                System.gc();
            } finally {
                System.out.println("System.gc() called");
                release();
            }
        } catch (Exception ex) {
            // ignore the InterruptedException
        }
    }
}
```

Calling System.gc() is bad (but not **the** problem)

Empty catch block hides problem

Lab 4: Combining Your Skills

Avoiding Liveness Hazards



Lab 4: Combining your skills

- **Problem: try to solve lab 4 using the skills learned**
- **Be careful - it is not as easy as it looks :-)**

Lab 5: Speeding Up Fibonacci



Lab 5: Speeding Up Fibonacci

- **Number sequence named after Leonardo of Pisa**

- $F_0 = 0$

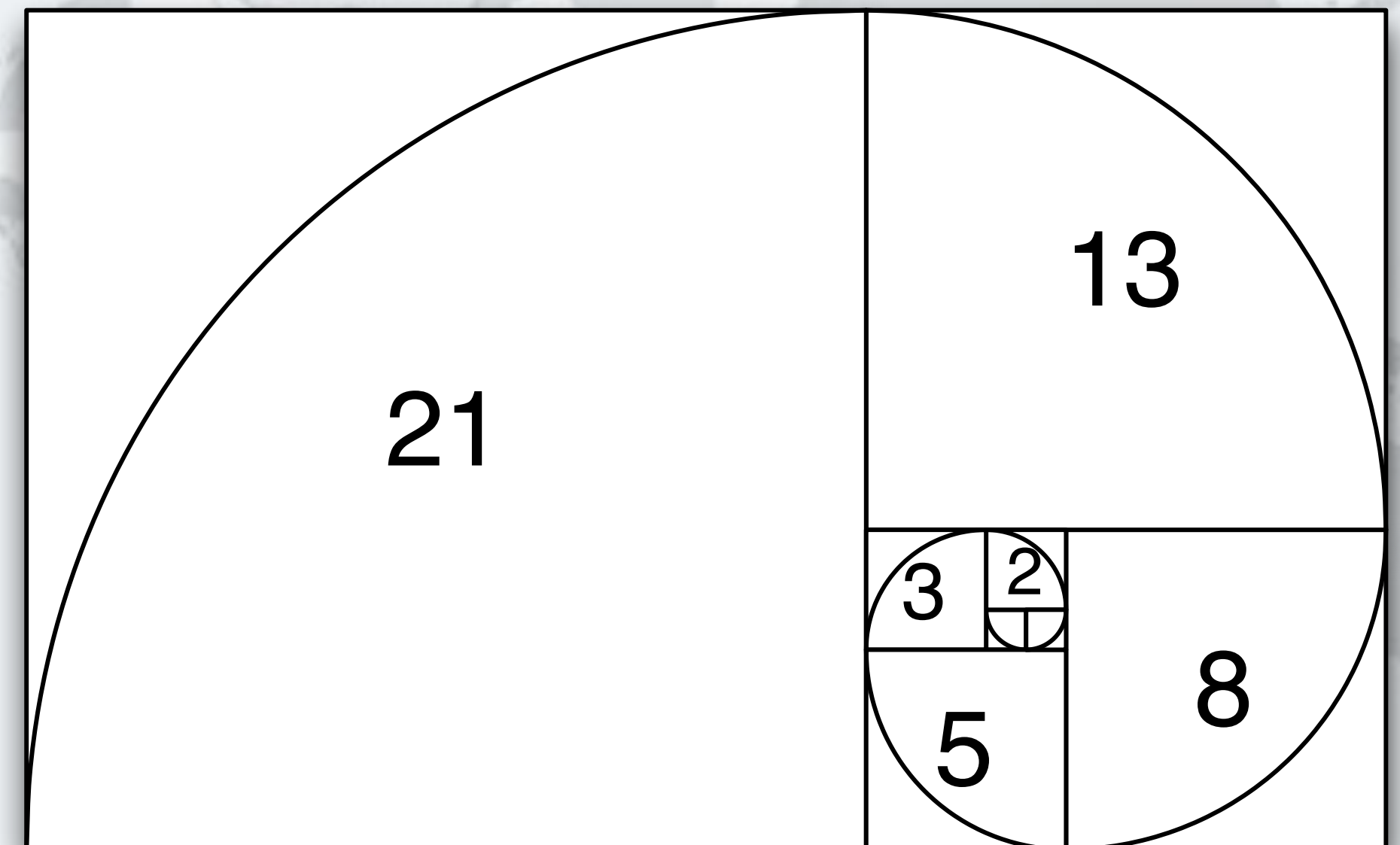
- $F_1 = 1$

- $F_n = F_{n-1} + F_{n-2}$

- **Thus the next number is equal to the sum of the two previous numbers**

- e.g. 0, 1, 1, 2, 3, 5, 8, 13, 21, ...

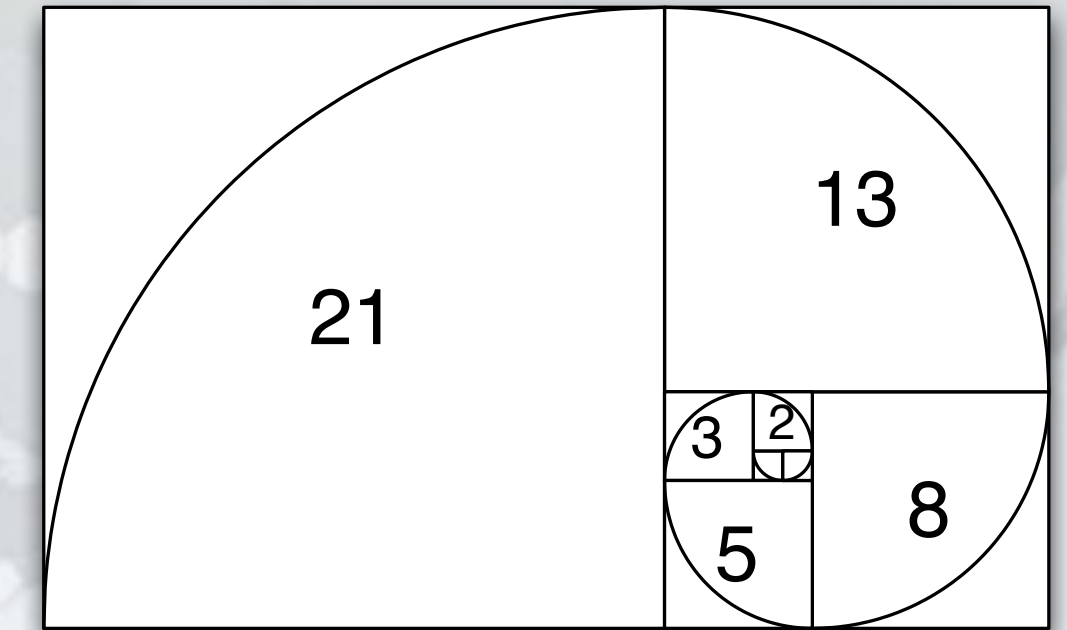
- **The numbers get large very quickly**



Exponential Algorithm

- **Taking our recursive definition**

- $F_0 = 0, F_1 = 1$
- $F_n = F_{n-1} + F_{n-2}$



- **Our first attempt writes a basic recursive function**

```
public long f(int n) {  
    if (n <= 1) return n;  
    return f(n-1) + f(n-2);  
}
```

- **But this has exponential time complexity**

- $f(n+10)$ is 1000 slower than $f(n)$

Linear Algorithm

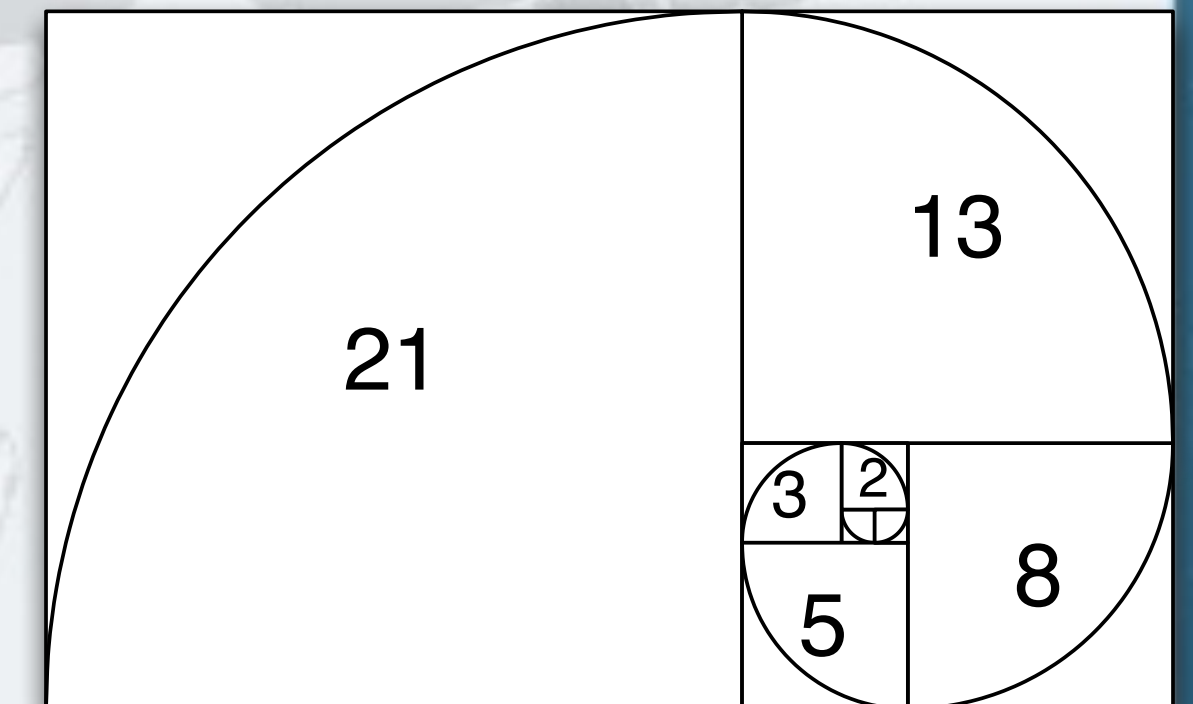
- **Instead of a recursive method, we could use iteration:**

```
public static long f(int n) {  
    long n0 = 0, n1 = 1;  
    for (int i = 0; i < n; i++) {  
        long temp = n1;  
        n1 = n1 + n0;  
        n0 = temp;  
    }  
    return n0;  
}
```

- **This algorithm has linear time complexity**

- **Solved $f(1_000_000_000)$ in 1.7 seconds**

- **However, the numbers overflow so the result is incorrect**
- **We can use `BigInteger`, but its `add()` is also linear, so time is quadratic**
- **We need a better algorithm**

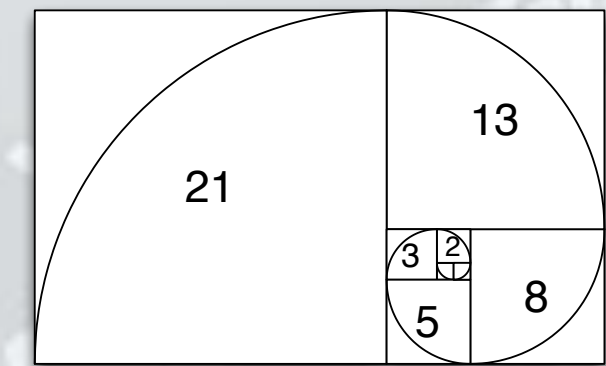


3rd Attempt Dijkstra's Sum of Squares

- **Dijkstra noted the following formula for Fibonacci**

- $F_{2n-1} = F_{n-1}^2 + F_n^2$

- $F_{2n} = (2 \times F_{n-1} + F_n) \times F_n$



- **Logarithmic time complexity and can be parallelized**

- Java 8 uses better BigInteger multiply() algorithms

- Karatsuba complexity is $O(n^{1.585})$

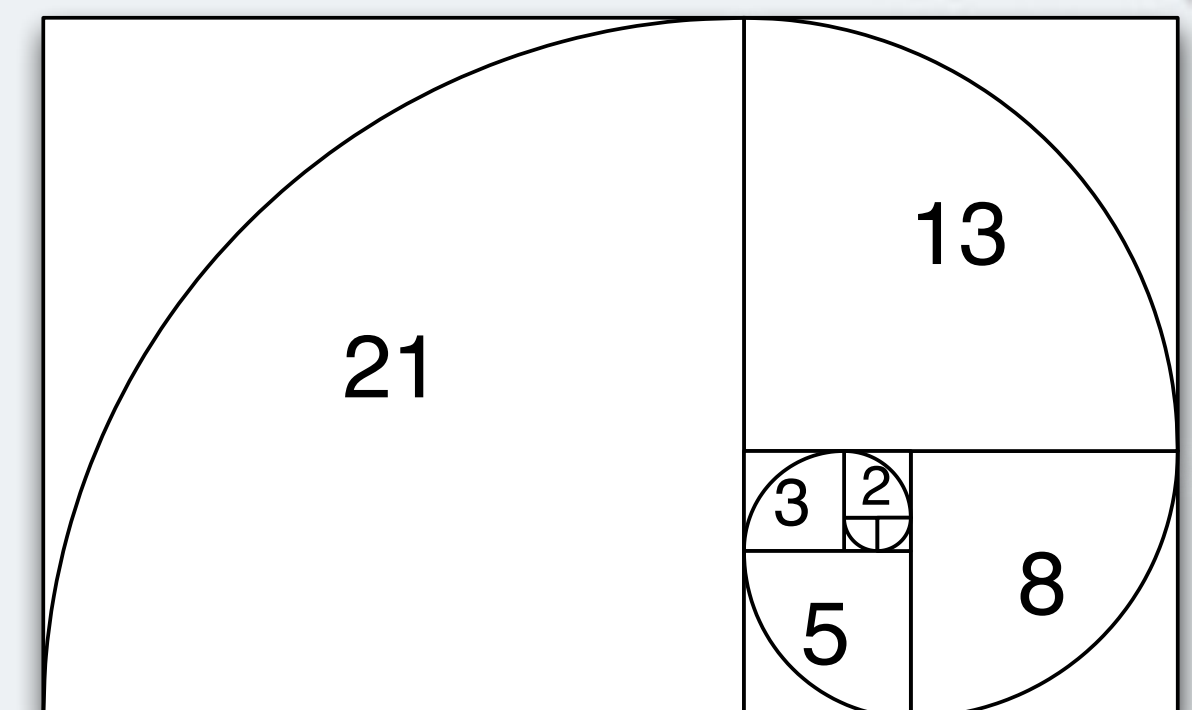
- 3-way Toom Cook complexity is $O(n^{1.465})$

- Previous versions of Java had complexity $O(n^2)$

- Single-threaded - we'll fix that in Lab 5.3

Lab 5.1: Dijkstra's Sum of squares

- **Implement this algorithm using BigInteger**
 - $F_{2n-1} = F_{n-1}^2 + F_n^2$
 - $F_{2n} = (2 \times F_{n-1} + F_n) \times F_n$
- **Run all tests in FibonacciTest and record the times**
- **Do it yourself - no cheating with Google!**



Lab 5.2: Parallelize your algorithm

- **We can parallelize by using common Fork/Join Pool**
 - Next we fork() the 1st task, do the 2nd and then join 1st

```
ForkJoinTask<BigInteger> f0_task = new RecursiveTask<BigInteger>() {  
    protected BigInteger compute() {  
        return f(half - 1);  
    }  
}.fork();  
BigInteger f1 = f(half);  
BigInteger f0 = f0_task.join();
```


Lab 5.3: Parallelize BigInteger

- **Using principles from lab 5.2, parallelize methods in `eu.javaspecialists.performance.math.BigInteger`**
 - `multiplyToomCook3()`
 - `squareToomCook3()`
- **These would probably not reach the threshold, so we won't parallelize them:**
 - `multiplyKaratsuba()`
 - `squareKaratsuba()`

Lab 5.4: Cache Results

- **Dijkstra's Sum of Squares needs to work out some values several times. Cache results to avoid this.**
- **Make sure you implement a “reserved caching scheme” where if one thread says he wants to calculate some value, others would wait**
 - **e.g. have a special BigInteger that signifies RESERVED**
 - **First thing a task would do is check if map contains that**
 - **If it doesn't, it puts it in and thus reserves it**
 - **If it does, it waits until the task is done and uses that value**

Lab 5.5: ManagedBlocker

- **ForkJoinPool is configured with *desired parallelism***
 - Number of active threads
 - ForkJoinPool mostly used with CPU intensive tasks
- **If one of the FJ Threads has to block, a new thread can be started to take its place**
 - This is done with the ManagedBlocker
- **Change your cache to use ManagedBlocker to keep parallelism high**

Wrap Up

Avoiding Liveness Hazards



Conclusion on Live(li)ness

- **Concurrency is difficult, but good tools and techniques solve problems**
- **These are just a few that we use**
- **For more, sign up to**
 - **The Java Specialists' Newsletter**
 - **tinyurl.com/jfokus2017**
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Finding and Solving Java Deadlocks

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